Name _____

_____ ID No. _____

There are 125 points total. (At 5 pts. = 1%, the first exam is 20% and this is 25%.)

1. (40 pts.) Indicate whether true or false. Beware of guessing:

correct answer +5 pts. incorrect answer -3 pts. no answer 0 pts

- (a) _____ Every finite set of strings is a regular language.
- (b) _____ If L is a Turing-recognizable language, so is \overline{L} .
- (c) _____ There are CFLs that Turing machines cannot recognize.
- (d) _____ A nondeterministic Turing machine can recognize more languages than a standard Turing machine can.
- (e) ____ A 2-stack PDA can recognize more languages than a standard 1-stack PDA can.
- (f) ____ A 2-tape Turing machine can recognize more languages than a standard 1-tape Turing machine can.
- (g) ____ There exists a Turing machine that can decide if two context free grammars generate the same language.
- (h) ____ There exists a Turing machine which can decide if two DFAs recognize the same language.
- 2. (25 pts.) Prove: If L is decidable, then $L^{\mathcal{R}}$ is decidable *Hint*: Make use of the Turing machine that decides L.

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- 3. (36 pts.) Give an example of a language which satisfies each of the following. If it is in the text (including Exercises and Problems), or is a simple modification of one of these, no proof is needed. Otherwise, give a proof.
 - (a) A CFL which is NOT regular.

(b) A decidable language which is NOT a CFL.

(c) A Turing-recognizable language which is NOT decidable.

- 4. (24 pts.) Let $L = \{ww^{\mathcal{R}} | w \in \{0, 1\}^*\}$. Construct either
 - (i) a context free grammar to generate L or
 - (ii) a PDA to recognize L.