

Sampling Random Variables

Bo Li, Math, UCSD, Spring 2019

Let (Ω, \mathcal{A}, P) be a probability space. Let $X: \Omega \rightarrow \mathbb{R}$ be a random variable (RV) with probability density function (PDF), or simply, the density $f_X: \mathbb{R} \rightarrow [0, \infty)$. Sampling X with the given density f_X means to produce X_1, X_2, \dots i.i.d (independent identically distributed) RVs with the common density/distribution f_X .

Often, a starting point is to generate $U \sim U[0,1]$ i.e., RV U that is uniformly distributed on $[0,1]$. This is a random number $\in [0,1]$. Then, using some method to generate/produce the needed RVs X_1, X_2, \dots i.i.d. with the given/target density f_X . $U \sim U[0,1] \Rightarrow X \sim f$.

Notation. $X \sim f$ means X is a RV distributed according to f , or with the PDF f .

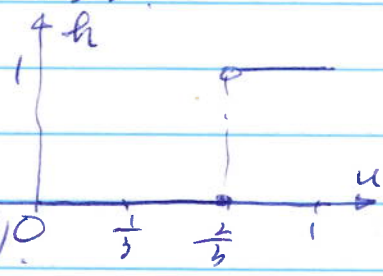
e.g., $U \sim U[0,1]$: U is uniformly distributed on $[0,1]$.
 $Z \sim N(0,1)$: Z has the standard normal or Gaussian distribution.

If $F = F(x)$ is the distribution (or cumulative distribution function - CDF) of a RV, then $X \sim F$ means the RV X has the CDF F , or X is F -distributed.

Example 1 Sample a RV $X \sim \text{Bernoulli}(\frac{1}{3})$

Define $h: [0, 1] \rightarrow \{0, 1\}$ by

$$h(u) = \begin{cases} 0 & \text{if } 0 \leq u \leq \frac{2}{3} \\ 1 & \text{if } \frac{2}{3} < u \leq 1 \end{cases}$$

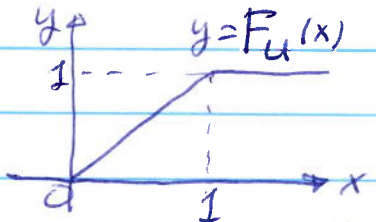


If $U \sim \mathcal{U}[0, 1]$ then $h(U) \sim \text{Bernoulli}(\frac{1}{3})$

Note that $P(U \notin [0, 1]) = 0$
 So we can define $h(U(\omega)) = 0$ if $U(\omega) \notin [0, 1]$.
 So $P(h(U) \notin \{0, 1\}) = 0$

Proof. $\forall \omega \in \Omega$.

$$h(U(\omega)) = \begin{cases} 0 & \text{if } U(\omega) \in [0, \frac{2}{3}] \\ 1 & \text{if } U(\omega) \in (\frac{2}{3}, 1] \end{cases}$$



So, $P(h(U) = 0) = P(0 \leq U \leq \frac{2}{3}) = F_U(\frac{2}{3}) - F_U(0) = \frac{2}{3}$

$P(h(U) = 1) = P(\frac{2}{3} < U \leq 1) = F_U(1) - F_U(\frac{2}{3}) = 1 - \frac{2}{3} = \frac{1}{3}$

Hence, $h(U) \sim \text{Bernoulli}(\frac{1}{3})$

Example 2 Sample $X \sim \text{Exp}(7)$. Recall the PDF

(or just the distribution)

for $\text{Exp}(7)$ [exponential distribution with $\lambda = 7$]

is given by $f(x) = \begin{cases} 0 & \text{if } x < 0 \\ 7e^{-7x} & \text{if } x \geq 0 \end{cases}$

Define $h: [0, 1] \rightarrow [0, \infty]$ by $h(u) = -\frac{1}{7} \log u$.

(log: natural logarithm).

Let $U \sim \mathcal{U}[0, 1]$. We show that $h(U) \sim \text{Exp}(7)$.

We compute the PDF (prob. distribution function) of the RV $X = h(U)$.

~~Since $h(u) \geq 0 \forall u \in [0, 1]$ we have $P(h(U) \leq 0) = 0$ i.e. $F_{h(U)}(x) = 0$ if $x \leq 0$~~

Let $t > 0$ we have

Note:

We should define

$$X = \begin{cases} h(U) & \text{if } U \in [0, 1] \\ 0 & \text{if } U \notin [0, 1] \end{cases}$$

But $P(U \leq 0 \text{ or } U > 1) = 0$

So, $F_X(x) = 0$ if $x \leq 0$

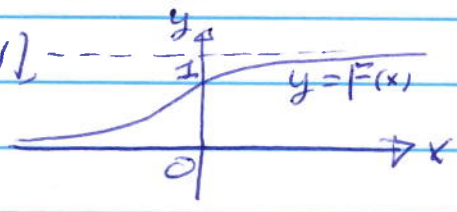
Let $t > 0$. We have

$$\begin{aligned}
 P(h(U) \leq t) &= P(\log U \geq -7t) \\
 &= P(U \geq e^{-7t}) \\
 &= P(U \geq e^{-7t}) + P(U < e^{-7t}) - P(U < e^{-7t}) \\
 &= 1 - P(U < e^{-7t}) \\
 &= 1 - F_U(e^{-7t}) \\
 &= 1 - e^{-7t} \quad (\text{since } e^{-7t} \in (0,1))
 \end{aligned}$$

i.e., $F_{h(U)}(x) = 1 - e^{-7x}$ if $x > 0$.
 Together $F_X(x) = \begin{cases} 0 & \text{if } x \leq 0 \\ 1 - e^{-7x} & \text{if } x > 0 \end{cases}$ $X = \begin{cases} h(U) & \text{if } 0 < U \leq 1 \\ 0 & \text{otherwise} \end{cases}$

The Inversion Method

Lemma Let $F: \mathbb{R} \rightarrow [0,1]$ be the distribution of a RV $: \Omega \rightarrow \mathbb{R}$. If $U \sim U[0,1]$, then $X := F^{-1}(U) \sim F$.



Note: We should really define $X = \begin{cases} F^{-1}(U) & \text{if } 0 \leq U \leq 1 \\ -\infty & \text{otherwise} \end{cases}$

Here F^{-1} is the generalized inverse of F , defined by

$$F^{-1}(u) = \inf \{ z \in \mathbb{R} : F(z) \geq u \} \quad \forall u \in (0,1)$$

If F is strictly increasing, then this is the same as the inverse function of F .

Prove this fact. You need the properties of the CDF F .

Proof of Lemma $\forall z \in \mathbb{R}$

$$\begin{aligned}
 P(F^{-1}(U) \leq z) &\stackrel{(*)}{=} P(U \leq F(z)) \\
 &\quad \uparrow \text{Use the right-continuity of } F \\
 &= F_U(F(z)) = F(z) \quad \text{since } F(z) \in [0,1]. \quad \square
 \end{aligned}$$

Pf of (*). Since $P(U \leq 0) = 0$, $P(U \geq 1) = 0$. We consider only $0 < U < 1$, and show $F^{-1}(U) \leq z \iff U \leq F(z)$ for any $z \in \mathbb{R}$. ~~\Leftarrow~~ Follows from the def. F^{-1} .

Then $P(X = -\infty) = P(U < 0 \text{ or } U > 1) = P(U < 0) + P(U > 1) = 0$.

\Rightarrow Fix $z \in \mathbb{R}$ and assume $F^{-1}(U) \leq z$. By the def. of $F^{-1}(U)$ \square
 $\exists z'_n \downarrow$ s.t. $F^{-1}(U) = \lim_{n \rightarrow \infty} z'_n$, and $F(z'_n) \geq U$. Note that z'_n is bounded below, for otherwise $z'_n \rightarrow -\infty$ and $F(z'_n) \rightarrow 0$ contradicting $F(z'_n) \geq U > 0$ ($n=1, 2, \dots$). So, $z'_n \downarrow z'$ for some $z' \in \mathbb{R}$. But all $z'_n \leq F^{-1}(U)$. So, $z' \leq F^{-1}(U) \leq z$. Moreover, the right-continuity implies $F(z') = \lim_{n \rightarrow \infty} F(z'_n) \geq U$. Hence, the monotonicity of F implies that $U \leq F(z') \leq F(z)$. \square

Example Sample a RV that is $\text{Exp}(\lambda)$ distributed for some $\lambda > 0$. The CDF for $\text{Exp}(\lambda)$ is $F(x) = \begin{cases} 0 & \text{if } x < 0 \\ 1 - e^{-\lambda x} & \text{if } x \geq 0 \end{cases}$. Let $U \sim U[0, 1]$. Since $P(U \notin [0, 1]) = 0$, let's just assume $U \in (0, 1)$. Let $X = F^{-1}(U)$.

$$x = F^{-1}(u), u \in (0, 1) \iff F(x) = u \in (0, 1)$$

$$1 - e^{-\lambda x} = u \in (0, 1), \iff x = -\frac{1}{\lambda} \log(1-u)$$

So, $X = -\frac{1}{\lambda} \log(1-U)$. Check that $F_X = F$.

So, $X \sim \text{Exp}(\lambda)$.

Example Sample a RV with density $f(r) = \begin{cases} 0 & \text{if } r < 0 \\ re^{-\frac{1}{2}r^2} & \text{if } r \geq 0 \end{cases}$. The distribution is $F(r) = \int_{-\infty}^r f(u) du = 1 - e^{-\frac{1}{2}r^2}$, if $r \geq 0$, and $F(r) = 0$ if $r \leq 0$. Solve $F(r) = u \in (0, 1)$

$$\Rightarrow r = \sqrt{-2 \log(1-u)}$$

So, if $U \sim U[0, 1]$. Then $R = \sqrt{-2 \log(1-U)}$ is a RV with the density f .

[See next page for the discrete inversion method.]

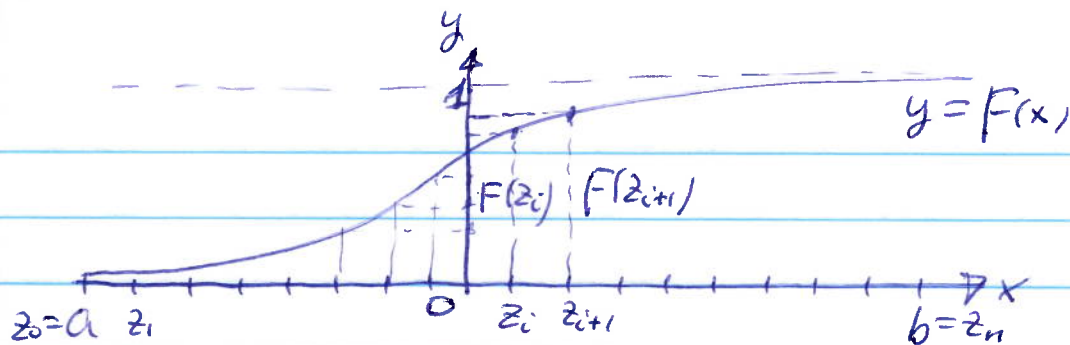
Numerical Approximation - if an explicit formula of F^{-1} is not available. Sample a RV with CDF F .

Step 1. Find $a, b \in \mathbb{R}$, $a < b$, and set

$$a = z_0 < z_1 < \dots < z_n = b.$$

So that $F(a) < \epsilon$, $1 - F(b) < \epsilon$, $|F(z_{i+1}) - F(z_i)| < \epsilon$ for all i .

Cont. from end of page 4.



Step 2 Generate $U \sim U[0,1]$.

Step 3 Find i s.t. $F(z_i) \leq U \leq F(z_{i+1})$
(Need to use some search algorithm)

Step 4. Define
$$X = z_i + (z_{i+1} - z_i) \frac{U - F(z_i)}{F(z_{i+1}) - F(z_i)}$$

Then, approximately, X has the distribution function F .

Why? By the inversion method, $F(X) = U$.
But, we don't have a formula for F^{-1} . So,
we solve $F(z) = U$ numerically to find z , then use this z as our X .

$$U = F(z) \approx F(z_i) + F'(z_i)(z - z_i)$$

$$\Rightarrow z \approx z_i + \frac{U - F(z_i)}{F'(z_i)} \quad \text{But } F'(z_i) = \frac{F(z_{i+1}) - F(z_i)}{z_{i+1} - z_i}$$

$$\text{Hence, } z \approx z_i + (z_{i+1} - z_i) \frac{U - F(z_i)}{F(z_{i+1}) - F(z_i)}$$

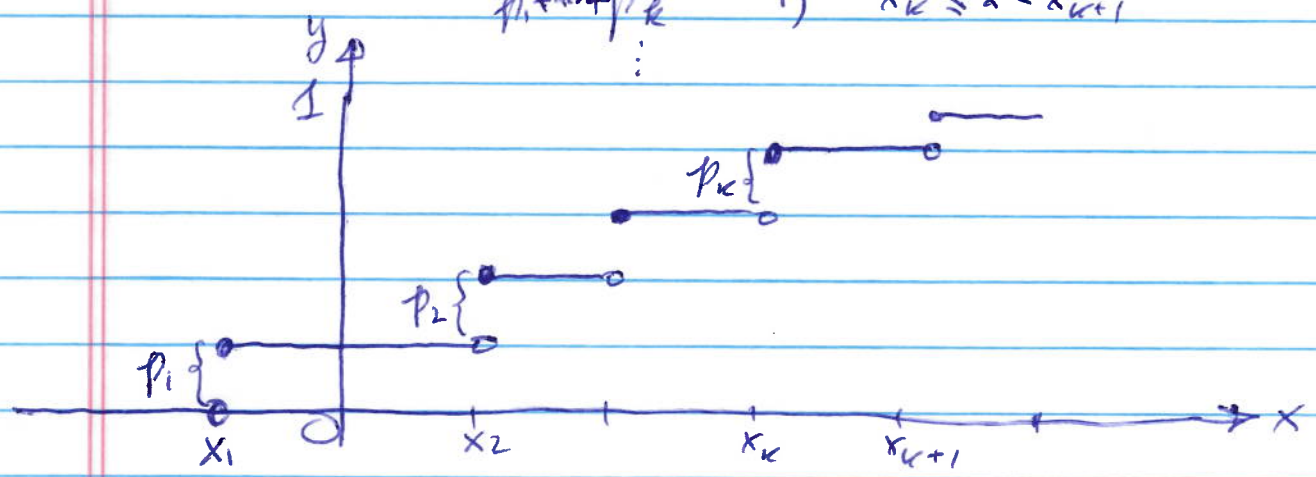
Replace z by X and " \approx " by " $=$ ".

The Discrete Inversion Method (Or the Inversion Method for Discrete RVs)

Let $X: \mathcal{R} \rightarrow \mathbb{R}$ be a discrete RV taking values $x_1 < x_2 < \dots$ with probabilities p_1, p_2, \dots , where $p_k = P(X = x_k) > 0$ and $\sum_{k \geq 1} p_k = 1$

The PMF (probability mass function), or the discrete distribution (function) $F = F_X: \mathbb{R} \rightarrow [0, 1]$ is a step function:

$$F(x) = \begin{cases} 0 & \text{if } x \leq x_1 \\ p_1 & \text{if } x_1 \leq x < x_2 \\ p_1 + p_2 & \text{if } x_2 \leq x < x_3 \\ \vdots & \vdots \\ p_1 + \dots + p_k & \text{if } x_k \leq x < x_{k+1} \end{cases}$$



If $U \sim U[0, 1]$. Then, for each $k \geq 1$, we find the smallest integer $k \geq 1$, such that $F(x_k) \geq U$, and return $X = x_k$. The X is a discrete RV with $X \sim F$.

This algorithm involves searching, which is in general $O(M)$ complexity if it is a M -state RV. With a binary search, it is $O(\log M)$.



The Transformation Method

More general than the inversion method.

Note:
 $\alpha = \alpha(x)$
may not
itself be
a distrib.
function.
 $\alpha(x)$ can
be negative

Let $X: \Omega \rightarrow \mathbb{R}$ be a RV with the distribution
 $F_X: \mathbb{R} \rightarrow [0, 1]$ and the density $f_X: \mathbb{R} \rightarrow [0, \infty)$
Let $\alpha: \mathbb{R} \rightarrow \mathbb{R}$ be a C^1 -function. Assume $\alpha'(x) > 0$
 $\forall x \in \mathbb{R}$. i.e. α is increasing.

Define $Y = \alpha(X)$. Then $Y: \Omega \rightarrow \mathbb{R}$ is
also a RV. We would like to find the
distribution and/or density for Y .

Let $y = \alpha(x)$ ($x \in \mathbb{R}$) Then

$$\begin{aligned} F_Y(y) &= P(Y \leq y) = P(\alpha(X) \leq \alpha(x)) \\ &= P(X \leq x) = F_X(x) \end{aligned}$$

$$\begin{aligned} F_Y'(y) &= f_Y \\ F_X'(x) &= f_X \end{aligned} \quad \text{So, } \frac{d}{dx} [F_Y(\alpha(x))] = f_X'(x) = f_X(x)$$

$$F_Y'(y) \alpha'(x)$$

$$F_Y(y) \frac{dx}{dy} = f_X(x)$$

$$\boxed{f_Y(y) = f_X(x) \left(\frac{dx}{dy} \right)^{-1}, \quad \text{given } x = \alpha^{-1}(y)}$$

Let $\beta \in C^1(\mathbb{R})$ $\beta \downarrow_n$ strictly i.e., $\beta'(x) < 0 \forall x \in \mathbb{R}$.

Consider now $Y = \beta(X)$.
What is f_Y ?

$$\begin{aligned} \text{Let } y = \beta(x). \quad F_Y(y) &= P(Y \leq y) = P(\beta(X) \leq y(x)) \\ &= P(X \geq x) = 1 - P(X < x) \\ &= 1 - F_X(x) \end{aligned}$$

$$\frac{d}{dy}: \quad f_Y(y) \frac{d\beta}{dx} = -f_X(x)$$

$$\boxed{f_Y(y) = f_X(x) \left| \frac{d\beta}{dx} \right|^{-1} \quad y = \beta^{-1}(x)}$$

For both cases $|f_X(x) dx| = |f_Y(y) dy|$.

Example $X \sim N(0,1)$ $f_X(x) = \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}}$

Let $\alpha(x) = \sigma x + \mu$, for some $\sigma > 0$ and $\mu \in \mathbb{R}$.

$$\text{Let } Y = \alpha(X). \quad f_Y(y) = f_X(x) \left(\frac{dy}{dx} \right)^{-1} = \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}} \cdot \frac{1}{\sigma}, \quad y = \alpha(x)$$

But $y = \alpha(x) = \sigma x + \mu$. So, $x = \frac{y - \mu}{\sigma}$. Hence,

$$f_Y(y) = \frac{1}{\sigma \sqrt{2\pi}} e^{-\frac{(y - \mu)^2}{2\sigma^2}}$$

And, $Y \sim N(\mu, \sigma^2)$.

Example $X \sim U[0,1]$ $\alpha(x) = -\log x, \quad x \in [0,1]$

Let $Y = \alpha(X) = -\log X$. For $y = \alpha(x), \quad x \in [0,1]$,

$$f_Y(y) = f_X(x) \left| \frac{d\alpha}{dx} \right|^{-1} = \left| \frac{-1}{x} \right|^{-1} f_X(x) = x f_X(x)$$

If $x \in (0,1)$ then $y = \alpha(x) = -\log x \in (0, \infty)$

Also, $X \sim U[0,1]$. So $f_X(x) = 1$ for $x \in (0,1)$

Hence $f_Y(y) = x = e^{-y}$ for $y \in (0, \infty)$.

Note that $P(X \leq 0) = F_U(0) = 0$
 $P(X \geq 1) = 0$ ($= 1 - P(X < 1) = 1 - 1 = 0$)

So, on the subset $\{X \leq 0\} \in \mathcal{R}$, we can define $Y = 0$. So,

$$Y = \begin{cases} 0 & \text{on } \{X \leq 0\} \\ -\log X & \text{on } \{X > 0\} \end{cases} \quad \{0 < X \leq 1\}$$

Thus,

$$F_Y(y) = P(Y \leq y) = 0 \text{ if } y \leq 0.$$

Hence $f_Y(y) = F'_Y(y) = 0$
if $y \leq 0$

Now, for $y > 0$

$$F'_Y(y) = \int_{-\infty}^y f_Y(u) du = e^{-y}$$

Hence,
$$F_Y(y) = \int_{-\infty}^y f_Y(u) du = \int_0^y e^{-u} du = 1 - e^{-y}$$

Summary:
$$F_Y(y) = \begin{cases} 0 & \text{if } y < 0 \\ 1 - e^{-y} & \text{if } y \geq 0 \end{cases}$$

This is the exponential distribution with $\lambda = 1$.

In what follows, if $U \sim U[0,1]$, i.e., $U: \mathcal{R} \rightarrow \mathcal{R}$ is a RV $U[0,1]$ -distributed. Then we can assume $0 \leq U \leq 1$. Since $P(U \notin [0,1]) = 0$.

In computer, a RND (a random number) generated in $[0,1]$ is always in $[0,1]$.

The transformation method for sampling multi-variate (multiple RVs).



Let $X_1, \dots, X_n: \Omega \rightarrow \mathbb{R}$ be n RVs. Then, $X = (X_1, \dots, X_n): \Omega \rightarrow \mathbb{R}^n$ is a vector-valued RV.

The joint distribution function is $F_X: \mathbb{R}^n \rightarrow [0, 1]$.

$$F_X(x) = P\left(\bigcap_{i=1}^n \{X_i \leq x_i\}\right) \quad \forall x = (x_1, \dots, x_n) \in \mathbb{R}^n.$$

The joint density (assumed to exist) is denoted

$$f_X(x) = f_{X_1, \dots, X_n}(x_1, \dots, x_n): \mathbb{R}^n \rightarrow [0, \infty)$$

It is related to F_X by

$$F_X(x) = \int_{-\infty}^{x_1} \dots \int_{-\infty}^{x_n} f(u_1, \dots, u_n) du_1 \dots du_n$$

Let $T: \mathbb{R}^n \rightarrow \mathbb{R}^n$ be a C^1 -map with the Jacobian for $T^{-1}: \text{Range}(T) \rightarrow \mathbb{R}^n$ ^{and invertible}

$$J(y) = \det\left(\frac{\partial x_i}{\partial y_j}\right)$$

nonsingular at any $y \in \text{Range}(T)$, where

$$T(x_1, \dots, x_n) = (y_1, \dots, y_n),$$

$$(x_1, \dots, x_n) = T^{-1}(y_1, \dots, y_n).$$

$$x_i = x_i(y_1, \dots, y_n) \quad (i=1, \dots, n).$$

Define $Y = T(X_1, \dots, X_n): \Omega \rightarrow \mathbb{R}^n$. What is the joint density for Y ?

Theorem Under the above assumptions,

we have for $Y = T(X_1, \dots, X_n) = (Y_1, \dots, Y_n)$,

$$f_Y(y) = \begin{cases} f_X(T^{-1}(y)) |J(y)| & \text{if } y \in \text{Range}(T), \\ 0 & \text{otherwise.} \end{cases}$$

PF Assume $\text{Rang}(T) = \mathbb{R}^n$. Write $x = (x_1, \dots, x_n)$,
 $y = (y_1, \dots, y_n)$. $\forall B \subseteq \mathbb{R}^n$ a Borel set. [It is
enough to assume B is a box: $B = \prod_{i=1}^n [a_i, b_i]$.]

We have

$$P(X \in B) = \int_B f_X(x) dx = \int_{T^{-1}(B)} f_X(T^{-1}(y)) |J(y)| dy$$

Since $Y = T(X)$: $P(Y \in T(B)) = \int_{T(B)} f_Y(y) dy$

Hence $f_Y(y) = f_X(T^{-1}(y)) |J(y)|$. \square

The result is $f_Y(y) dy = f_X(x) dx$. $\frac{dx}{dy} = |J(y)|$
formally:

The Box-Muller Method for Sampling Two Independent Gaussian RVs and standard indep.

Let $X, Y: \Omega \rightarrow \mathbb{R}$ be two standard Gaussian variables.

(1) Independent: $f_{X,Y}(x,y) = f_X(x) f_Y(y)$

(2) $f_X(x) = \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}}$ $f_Y(y) = \frac{1}{\sqrt{2\pi}} e^{-\frac{y^2}{2}}$

~~(2)~~ i.e. $X \sim \mathcal{N}(0,1)$, $Y \sim \mathcal{N}(0,1)$

We would like to sample X, Y by generating $z, \eta \sim U[0,1]$ and use the transformation method.

We implement this using two steps.

Step 1. Change variables (x, y) to (r, θ)

$$x = r \cos \theta, \quad X = R \cos \Theta$$

$$y = r \sin \theta, \quad Y = R \sin \Theta$$

Generate RVs R and Θ that correspond to X, Y .

Find f_R, f_Θ

Step 2. Generate $z, \eta \sim U[0, 1]$. And then using the transformation method again to generate R and Θ .

or the inversion method

Let's begin with

$$f_{X,Y}(x,y) dx dy = f_{R,\Theta}(r,\theta) dr d\theta$$

Independence:

$$f_X(x) f_Y(y) dx dy = f_R(r) f_\Theta(\theta) dr d\theta$$

But, the left-hand side is

$$\int_{x=r \cos \theta} \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}} \int_{y=r \sin \theta} \frac{1}{\sqrt{2\pi}} e^{-\frac{y^2}{2}} dx dy = \frac{1}{2\pi} e^{-\frac{1}{2}(x^2+y^2)} dx dy = \frac{1}{2\pi} e^{-\frac{1}{2}r^2} r dr d\theta$$

So, define R, Θ so that $X = R \cos \Theta$ $Y = R \sin \Theta$.

Hence, we choose $f_R(r) = r e^{-\frac{1}{2}r^2} \quad (r \geq 0)$
 $f_\Theta(\theta) = \frac{1}{2\pi} \quad \theta \in [0, 2\pi)$

We can define $f_R(r) = 0$ if $r < 0$ and $f_\Theta(\theta) = 0$ if $\theta \notin [0, 2\pi)$.

Now, (step 2) we generate R, Θ with the densities f_R and f_Θ , respectively, by the inversion method.

By the second example on page [4] of this set of notes, we have

$$R = \sqrt{-2 \log(1-\xi)} \quad \text{if } \xi \sim U[0,1]$$

For sampling Θ , note first from $f_{\Theta}(\theta) = \frac{1}{2\pi}$ only for $\theta \in [0, 2\pi)$, we have

$$F_{\Theta}(\alpha) = \begin{cases} 0 & \text{if } \alpha \leq 0 \\ \frac{\alpha}{2\pi} & \text{if } 0 < \alpha \leq 2\pi \\ 1 & \text{if } \alpha > 2\pi. \end{cases}$$

Let $\eta \sim U[0,1]$.

Solve: $F_{\Theta}(\alpha) = \eta$ in $0 < \alpha \leq 2\pi$.

Hence $\frac{\alpha}{2\pi} = \eta, \alpha = 2\pi \eta \in [0, 2\pi)$

$$R_{\Theta} = 2\pi \eta$$

Now, we have the Box-Muller method

$$\begin{aligned} X &= \sqrt{-2 \log(1-\xi)} \cos(2\pi \eta) \\ Y &= \sqrt{-2 \log(1-\xi)} \sin(2\pi \eta) \end{aligned} \quad \begin{matrix} \xi, \eta \sim U[0,1] \\ \text{independent} \end{matrix}$$

Let us verify that $X, Y \sim N(0,1)$, independent.

Since $\xi, \eta \sim U[0,1]$ independent,

$$R = \sqrt{-2 \log(1-\xi)}$$

$$\Theta = 2\pi \eta$$

are independent RVs. So, $f_{R, \Theta} = f_R \cdot f_{\Theta}$.

Since $\eta \sim U[0,1]$ \rightarrow

$$F_{\Theta}(\alpha) = P(\Theta \leq \alpha) = P(\eta \leq \frac{\alpha}{2\pi}) = F_{\eta}(\frac{\alpha}{2\pi})$$

$$= \begin{cases} \frac{\alpha}{2\pi} & \text{if } 0 \leq \alpha \leq 2\pi \\ 1 & \text{elsewhere if } \alpha > 2\pi. \end{cases}$$

So, $f_{\Theta}(\theta) = \begin{cases} \frac{1}{2\pi} & \text{if } \theta \in [0, 2\pi) \\ 0 & \text{elsewhere.} \end{cases}$

Similarly,

$$\begin{aligned} \forall r \geq 0 \quad F_R(r) &= P(R \leq r) = P(\xi \geq e^{-\frac{r^2}{2}}) \\ &= 1 - P(\xi < e^{-\frac{r^2}{2}}) \\ &= 1 - F_\xi(e^{-\frac{r^2}{2}}) \\ &= 1 - e^{-\frac{r^2}{2}} \quad \text{since } \xi \sim U([0,1]) \end{aligned}$$

$$\forall r < 0 \quad F_R(r) = P(R \leq r) = 0.$$

$$\text{So, } f_R(r) = \begin{cases} 0 & \text{if } r < 0 \\ F_R'(r) & \text{if } r \geq 0 \end{cases} = \begin{cases} 0 & \text{if } r < 0, \\ re^{-\frac{r^2}{2}} & \text{if } r \geq 0. \end{cases}$$

$$\text{Combine: } f_{R,\Theta}(r,\theta) = f_R(r) f_\Theta(\theta) = \begin{cases} \frac{1}{2\pi} re^{-\frac{r^2}{2}} & \text{if } r \geq 0, \theta \in [0, 2\pi] \\ 0 & \text{elsewhere.} \end{cases}$$

Now, by the theorem on the transformation method for multi-variables, we have for the new RVs $X = R \cos \Theta$, $Y = R \sin \Theta$. following the transform. $T: (R, \Theta) = (x, y)$
 $x = r \cos \theta, y = r \sin \theta.$

$$\text{that } f_{X,Y}(x,y) = f_{R,\Theta}(r,\theta) |J(x,y)|$$

$$|J(x,y)| = \left| \frac{\partial(r,\theta)}{\partial(x,y)} \right| = \left| \det \begin{pmatrix} \frac{\partial r}{\partial x} & \frac{\partial r}{\partial y} \\ \frac{\partial \theta}{\partial x} & \frac{\partial \theta}{\partial y} \end{pmatrix} \right| = \frac{1}{r}$$

for $r \geq 0$. Hence

$$\begin{aligned} f_{X,Y}(x,y) &= f_{R,\Theta}(r,\theta) \frac{1}{r} \\ &= \frac{1}{2\pi} e^{-\frac{r^2}{2}} \\ &= \frac{1}{2\pi} e^{-\frac{1}{2}(x^2+y^2)} \end{aligned}$$

The marginal distribution is

$$\begin{aligned} F_X(x) &= F_{X,Y}(x, \infty) = \int_{-\infty}^{\infty} dy \int_{-\infty}^x du f_{X,Y}(u,v) \\ &= \int_{-\infty}^x \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2}u^2} du = \left[\frac{1}{\sqrt{2\pi}} \int_{-\infty}^x e^{-\frac{1}{2}u^2} du = 1. \right] \end{aligned}$$

Hence $f_X(x) = \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}}$

Similarly $f_Y(y) = \frac{1}{\sqrt{2\pi}} e^{-\frac{y^2}{2}}$

Hence, $X, Y \sim N(0, 1)$. Moreover

$f_{X,Y}(x,y) = f_X(x) f_Y(y)$

So, X and Y are independent.

Note. To generate a single RV $X \sim N(0, 1)$, one can use the CLT (Central Limit Theorem)

(1) Generate $\xi_1, \xi_2, \dots, \xi_N$ i.i.d. $U[0, 1]$ -distributed.

(2) Set $X_N = \sqrt{2N} \left(\sum_{k=1}^N \xi_k - \frac{1}{2}N \right)$

Then, Approximately $X_N \sim N(0, 1)$. if $N \gg 1$.

Accept-Reject (or Acceptance-Rejection)

The Acceptance-Rejection Method

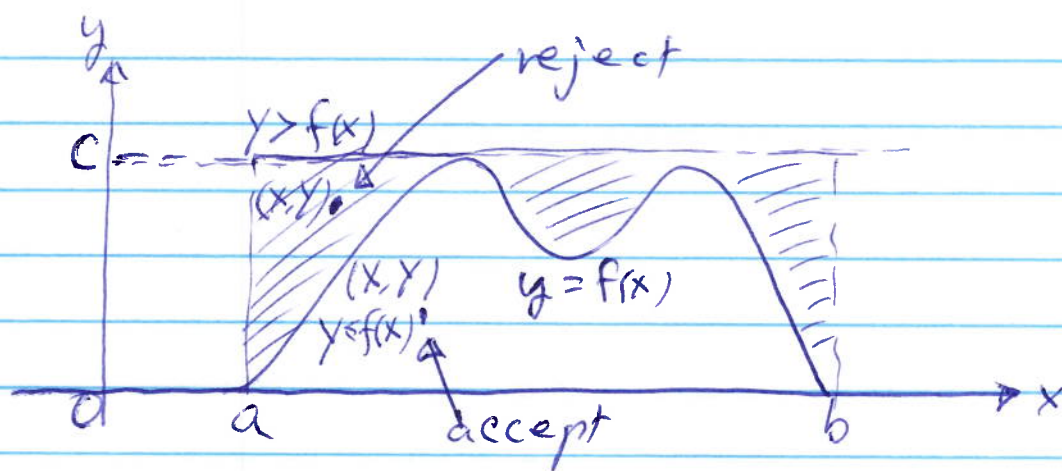
(Due to John von Neumann early 1950s.) necessary

transforming from

This is a direct sampling method, without using some other distributions (e.g., $U[0, 1]$). So, it is different from the inversion, or more generally, the transformation method.

Basic idea. Let $f: \mathbb{R} \rightarrow [0, \infty)$ be the pdf of a RV from \mathbb{R} to \mathbb{R} . Assume $f = 0$ outside $[a, b]$, a finite interval ($a, b \in \mathbb{R}, a < b$) We would like to sample a RV with this pdf.

(The assumption $f = 0$ outside $[a, b]$ can be relaxed, to get approximate results.)



Let $c \in \mathbb{R}$ be an upper bound of f , e.g.,
 (Simple acceptance-rejection) $c = \sup \{ f(x) : x \in [a, b] \}$.

The algorithm for generating $Z \sim f$ is:

Step 1 Generate $X \sim U[a, b]$

~~Step 1~~ e.g., generate $U \sim U[0, 1]$.

$$X \leftarrow (b-a)U + a$$

Step 2 Generate $Y \sim U[0, c]$, independent of X . accept X and

Step 3 If $Y \leq f(X)$, return $Z = X$.

Otherwise, reject X and return to step 1.

Note. Here, we still use a reference distribution $U[0, 1]$.

Why the algorithm generates $Z \sim f$? The reason is as follows. (X, Y) is uniformly distributed ^{on $[a, b] \times [0, c]$} , and the accepted X satisfies $Y \leq f(X)$. Thus the marginal distribution for the accepted X is f . and X and Y are independent.

marginal: $f_X(x) = \int_0^{f(x)} 1 dy = f(x)$.

A more general ^{accept-reject or} ~~acceptance-rejection~~ method 17
is based on the following idea:

Let $f: \mathbb{R}^n \rightarrow [0, \infty)$ be the PDF of a RV in \mathbb{R}^n . Let $g: \mathbb{R}^n \rightarrow (0, \infty)$ be the PDF of a RV in \mathbb{R}^n , and $\alpha \geq 1$ be such that

① We know how to sample RVs with g the PDF.

② $\alpha g(x) \geq f(x)$ for all x .

But the region between the surfaces $y = \alpha g(x)$ and $y = f(x)$ is as small as possible. So, $\alpha = \sup_x \frac{f(x)}{g(x)}$ is a good candidate.



Call $\alpha g(x)$ a majorizing function of $f(x)$, and $g(x)$ a proposal PDF.

Algorithm: Acceptance-Rejection Method for generating

step 1 Generate $X \sim g$.

step 2 Generate $Y \sim U[0, \alpha g(X)]$.

step 3 If $Y \leq f(X)$ then accept, $Z \leftarrow X$.
Otherwise, reject. Return to step 1.

Theorem The algorithm generates $Z \sim f$.

Proof Let $A = \{(x, y) : 0 \leq y \leq \alpha g(x)\} \subseteq \mathbb{R}^n \times \mathbb{R}^1$, □
 $B = \{(x, y) : 0 \leq y \leq f(x)\} \subseteq \mathbb{R}^n \times \mathbb{R}^1$. □

Note that the volume of A is $|A| = \alpha$, since g is a PDF, and the volume of B is $|B| = 1$, since f is a PDF.

From steps 1 and 2, (X, Y) is uniformly distributed on A . To see this let $q(x, y)$ denote the joint PDF of (X, Y) , and let $q(y|x)$ denote the conditional PDF of Y given $X=x$. Then,

$$q(x, y) = \begin{cases} g(x) q(y|x) & \text{if } (x, y) \in A, \\ 0 & \text{otherwise.} \end{cases}$$

By step 2 implies that $q(y|x) = \frac{1}{\alpha g(x)}$ for $y \in [0, \alpha g(x)]$, and $q(y|x) = 0$ elsewhere. Thus,

$$q(x, y) = \frac{1}{\alpha} \text{ for every } (x, y) \in A.$$

Let (X^*, Y^*) be the first accepted point, i.e., the first point in B . Since (X, Y) is uniformly distributed on A , (X^*, Y^*) is uniformly distributed on B . But, the volume $|B| = 1$. So, this joint PDF of (X^*, Y^*) is 1. Thus, the marginal PDF of $Z = X^*$ is

$$\int_0^{f(x)} 1 dy = f(x). \quad \square$$

Note that the efficiency of this algorithm is defined as

$$P((X, Y) \text{ is accepted}) = \frac{\text{vol.}(B)}{\text{vol.}(A)} = \frac{1}{\alpha}.$$

Also, in generating many RVs $\sim f$, we introduce a similar concept:

$$\text{acceptance rate} = \frac{\# \text{ of } X \text{ accepted}}{\# \text{ of } X \text{ generated}} \quad (= \text{efficiency})$$

Often a slightly modified version of the above algorithm is used.

$Y \sim U[0, \alpha g(X)]$ is same as setting $Y = U \alpha g(X)$, where $U \sim U[0, 1]$.

Then, $Y \leq f(X)$ is equivalent to $U \leq \frac{f(X)}{\alpha g(X)}$.

Hence, the modification is:

Generate X from $g(x)$ and accept it with probability $f(X)/[\alpha g(X)]$.

Algorithm: Modified Accept-Reject Method to

Sample a RV $Z \sim f$.

Let f be a given PDF of RV in \mathbb{R}^n .

Let g be a PDF of RV in \mathbb{R}^n and $\alpha \geq 1$

Suppose $f(x) \leq \alpha g(x) \forall x \in \mathbb{R}^n$ (known)

\mathbb{R}^n can be replaced by a finite region, the support of f , if f is finitely supported

Step 1 Generate X from g

Step 2 Generate $U \sim U[0, 1]$ independent of X . accept X , and

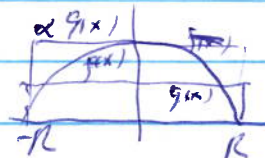
Step 3 If $U \leq \frac{f(X)}{\alpha g(X)}$, then $Z \leftarrow X$. Otherwise, reject X , ~~and~~ go to Step 1.

Example Let $f(x) = \frac{2}{\pi R^2} \sqrt{R^2 - x^2}$, $-R \leq x \leq R$, for some $R > 0$. $f \geq 0$ and $\int_{-R}^R f(x) dx = 1$. f is a PDF.

Take the proposed distribution: $g(x) = \frac{1}{2R}$ for $x \in [-R, R]$. Choose $\alpha = \text{const. small, s.t.}$

$$\alpha g(x) \geq f(x), \quad x \in [-R, R]$$

The smallest such α is $\alpha = \frac{7}{\pi}$.

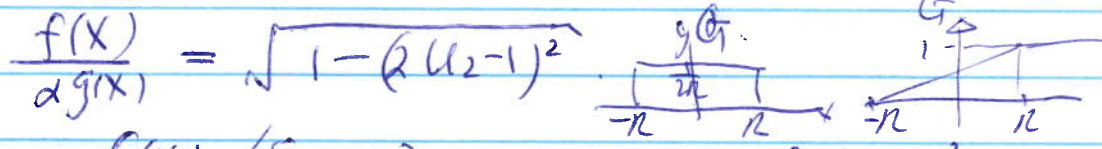


We sample a RV_Z with distribution f as follows.

① Generate independent $U_1, U_2 \sim U[0,1]$.

② Use U_2 to generate $X \sim g$ via the inversion method. $X = (2U_2 - 1)R$.

③ Calculate $\frac{f(X)}{\alpha g(X)} = \sqrt{1 - (2U_2 - 1)^2}$. ↖ Need first to calculate G



④ If $U_1 \leq \frac{f(X)}{\alpha g(X)}$, i.e.; $(2U_2 - 1)^2 \leq 1 - U_1^2$ or $U_1^2 + 4U_2^2 - 4U_2 \leq 0$, then accept X and return $Z = X = (2U_2 - 1)R$.

Otherwise, reject X and go to ①.

The expected number of trials is $\alpha = \frac{4}{\pi}$, and the efficiency is $\frac{1}{\alpha} = \frac{\pi}{4} \approx 0.785$.

~~Another~~ Another variation of the above algorithm is the so-called squeeze accept-reject method.

Let f be a PDF.

Let g, h also be PDFs that are relatively easy to sample. (e.g., $h(x)$ can be a piecewise linear function). Assume

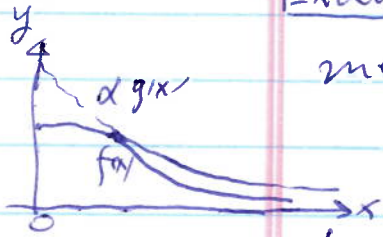
$$h(x) \leq f(x) \leq \alpha g(x) \quad \forall x.$$

Algorithm

Step 1 Generate $X \sim g(x), U \sim U[0,1]$.

Step 2 Accept X if $U \leq \frac{h(X)}{\alpha g(X)}$.

Step 3 Otherwise, accept X if $U \leq \frac{f(X)}{\alpha g(X)}$.



Example Sample $X \sim N(0,1)$ by the accept-reject method. The PDF is $f(x) = \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}}$ ($x \in \mathbb{R}$).

First, we choose a proposal PDF

$$g(x) = \frac{1}{2} e^{-|x|} \quad (x \in \mathbb{R}).$$

Observe that we can describe g as the distribution of an exponential ~~RV~~ RV ($\text{Exp}(\lambda)$, with $\lambda=1$) multiplied randomly by ± 1 . Thus, we can generate a RV $V \sim g$ as follows:

- (1) Generate $U_1, U_2 \sim U[0,1]$, independent,
- (2) If $U_1 \leq \frac{1}{2}$ then set $V = -\log U_2$.
Otherwise, set $V = \log U_2$.

Now, set

$$\alpha = \sup_{x \in \mathbb{R}} \frac{f(x)}{g(x)} = \sup_{x \geq 0} \sqrt{\frac{2}{\pi}} e^{x - \frac{x^2}{2}} = \sqrt{\frac{2e}{\pi}} \approx 1.32$$

So, $\alpha g(x)$ is a majorizing function of $f(x)$. i.e.,
 $0 \leq f(x) \leq \alpha g(x)$ ($x \in \mathbb{R}$)

The accept-reject method for generating $X \sim N(0,1)$

Step 1 Generate $V \sim g$ using $U_1, U_2 \sim U[0,1]$ independent, as in (1), (2) above.

Step 2 Generate $Y \sim U[\log g(V)]$ by generating $U_3 \sim U[0,1]$ and setting $Y = \log g(V) + U_3$
 $= \log \frac{1}{2} U_3 e^{-|V|} \quad (\Leftrightarrow U_3 \leq \frac{f(V)}{\alpha g(V)})$

Step 3 If $Y \leq f(V)$ then accept and set $X \leftarrow V$. Otherwise, reject and go to Step 1.

Efficiency. ~~Require~~ Require an average of $\alpha = 1.32$ proposals/trials (since acceptance rate $= 1/\alpha$). So, need ~~to~~ an average of $3 \times 1.32 = 3.96$ uniform random numbers to generate $X \sim N(0,1)$.

Discrete Accept-Reject Method

Generate a discrete RV X from a known target PMF $p: S \rightarrow [0,1]$. $p_i = P(X=x_i)$, $i \in S$.
 S is finite or countably infinite. $x_i \in \mathbb{R}$, distinct.

$$\underline{p_i \geq 0, \sum_{i \in S} p_i = 1.}$$

Suppose we know how to generate a RV V from a proposal PMF $q: S \rightarrow [0,1]$.

$q_i \geq 0$ ($i \in S$), $\sum_{i \in S} q_i = 1$. Suppose also there exists $\alpha \in \mathbb{R}$, $\alpha > 0$, such

$$p_i \leq \alpha q_i \quad \forall i \in S.$$

Then, we can proceed as follows:

- (1) Generate $V \sim q$.
- (2) Generate $Y \sim U[0, \alpha q]$.
- (3) If $Y \leq p_V$ then accept. Set $X \leftarrow V$ and stop. Otherwise, reject. Go to (1).

Example $S = \{1, 2, \dots\}$. The target PMF p is given by $p_i = \frac{6}{\pi^2 i^2}$ ($i=1, 2, \dots$). ($p_i \geq 0 \forall i$, $\sum_i p_i = 1$).

Take the proposal PMF to be q with $q_i = \frac{1}{i(i+1)}$ ($i=1, 2, \dots$). We can check that,

if $U \sim U[0,1]$, then $P(\lfloor U^{-1} \rfloor = i) = \frac{1}{i} - \frac{1}{i+1} = q_i$.

($\lfloor x \rfloor =$ greatest integer $\leq x$.) So, easy to generate

$V \sim q$. Also, $\max_{i \geq 1} \frac{p_i}{q_i} = \frac{p_1}{q_1} = \frac{12}{\pi^2}$.

Now, the algorithm generating $X \sim p$ is:

- step 1 Generate $U_1 \sim U[0,1]$ and set $V \leftarrow \lfloor U_1^{-1} \rfloor$.
- step 2 Generate $U_2 \sim U[0,1]$ and set $Y \leftarrow U_2 \left(\frac{12}{\pi^2} \frac{1}{V(V+1)} \right)$.
- step 3 If $Y \leq \frac{6}{\pi^2} V^{-2}$ then accept, set $X \leftarrow V$ and stop. Otherwise, Goto step 1.

The Alias Method for Generate Discrete RVs with Finite States. (Walker, 1974, 1977)

Let x_1, \dots, x_M be M distinct real numbers.

Denote $V = \{x_1, \dots, x_M\}$ and $S = \{1, 2, \dots, M\}$.

Let $X: \Omega \rightarrow V$ be a discrete RV such that

$$p_i := P(X = x_i) > 0, \quad \forall i \in S$$

and $\sum_{i=1}^M p_i = 1.$

So, the discrete density for X is given by

$$f(x) = \begin{cases} p_i & \text{if } x = x_i \text{ for some } i \in S. \\ 0 & \text{otherwise.} \end{cases}$$

We wish to sample ~~$Z \sim f$~~
a RV

The discrete inversion method: Generate $U \sim U[0,1]$. search m s.t. $p_1 + \dots + p_{m-1} < U \leq p_1 + \dots + p_{m-1} + p_m$ and set $Z \leftarrow x_m$. This involves the search of m with the complexity $O(M)$ for linear search or $O(\log M)$ for binary search.

The alias method consists of two parts, set up and sampling. Set up means the construction of M evenly weighted 2-point densities, equivalent to the original density, requiring work with the complexity same as that of the discrete inversion method. Then, sampling is to sample first M even values uniformly, followed by sampling a two-point density, i.e., generating a RV to be some x with ~~prob~~ some probability $p \in (0,1)$ and some y with probability $1-p$.

Set up: original distribution/density

p_1	p_2	...	p_M
1	2	...	M

$p_i > 0, \sum_{i=1}^M p_i = 1$

Construct $2 \times M$ table

Rules

5

If $p_k \leq \frac{1}{M}$, then p_k is one of a_j 's.
 If $p_k > \frac{1}{M}$, then p_k is splitted, into a few b_j 's and a possible a_j .

a_1	a_2	...	a_M
i_1	i_2	...	i_M
b_1	b_2	...	b_M
j_1	j_2	...	j_M

- ① $a_i \geq 0, b_i \geq 0$
- ② $i_j, j_k \in \{1, \dots, M\}$ (labels).
- ③ $a_k + b_k = \frac{1}{M}, k=1, \dots, M.$
- ④ $\sum_{k: i_k=r} a_k + \sum_{k: j_k=r} b_k = p_k, \forall r \in \{1, \dots, M\}$

Example $p_1=0.41, p_2=0.27, p_3=0.07, p_4=0.14, p_5=0.11$

$M=5$
 $\frac{1}{M} = 0.2$

0.41	0.27	0.07	0.14	0.11
1	2	3	4	5

Set up:

"Poor guys" = a_k 's with $a_k \leq \frac{1}{M}$, stand in 1st row.
 "Rich guys" = a_k 's with $a_k > \frac{1}{M}$. First, donates, then becomes "poor" ($\leq \frac{1}{M}$) then stand in Row 1.

0.07	0.11	0.14	0.19	0.20
3	5	4	1	2
0.13	0.09	0.06	0.01	0
1	1	2	2	2

Note: From M to $2M$, But, uniform in columns. $p_k = \frac{1}{M}$ = call it a middle "class" number
 $a_k + b_k = \frac{1}{M}, (k=1, 2, \dots, M)$

Note: It is possible that a rich p_k becomes 0 after donating its values to poor of numbers/guys. In this case, other rich numbers need to be two or more a_k values., e.g. 0.25, 0.06, 0.29, 0.3, 0.3.

Alias Algorithm

Set up (Construct M evenly weighted two-point densities, that are equivalent to the given density.)

Step 1 Set $S \leftarrow \{1, 2, \dots, M\}$.
Set $t \leftarrow 1$.
Set $r_k \leftarrow p_k$ ($k=1, 2, \dots, M$) for each $k \in S$.

Step 2 Set $i_t \leftarrow$ value of $k \in S$ that minimizes r_k .
Set $j_t \leftarrow$ value of $k \in S$ that maximizes r_k .
Set $a_t = r_{i_t}$.
Set $b_t = 1/M - a_t$.
Set $r_{j_t} \leftarrow r_{j_t} - b_t$.
Remove i_t from S .

Step 3 Set $t \leftarrow t+1$.
If $t \leq M$, go to step 2; otherwise stop.

Alias Algorithm

Sampling

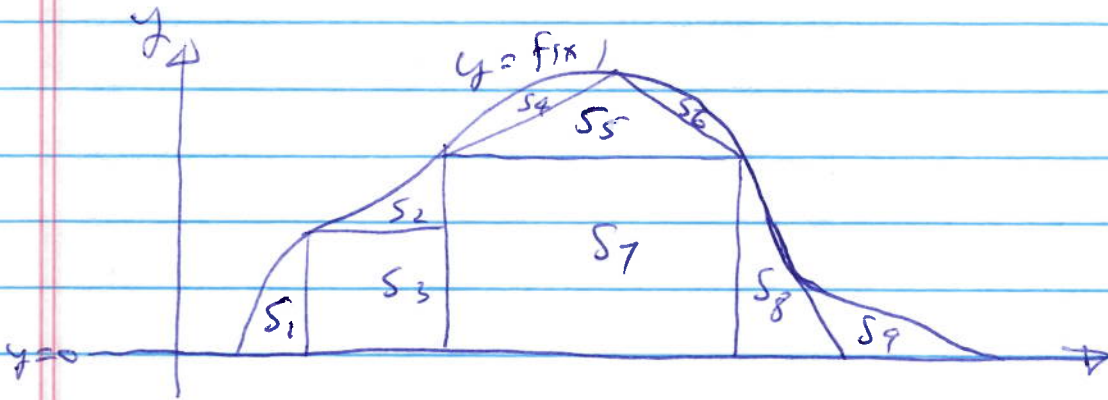
Step 1 Generate $U_1 \sim U[0,1]$ and set $K = \lceil MU_1 \rceil$.
($\lceil x \rceil =$ smallest integer $\geq x$.)

Step 2 Generate $U_2 \sim U[0,1]$. Set $Z \leftarrow X_{iK}$ if $U_2 \leq a_{iK}$. Otherwise set $Z \leftarrow X_{jK}$.

Theorem The algorithm generates $Z \sim f$.

Key Equivalence: $p_k = \frac{1}{M} \sum_{m=1}^M (a_m I_{[i_m=k]} + b_m I_{[j_m=k]}) \forall k \in S$.
PF this by induction.

The composition method



Let f be the PDF of a RV $X: \mathcal{R} \rightarrow \mathbb{R}$. $f_0, f \geq 0$
 $\int_{\mathcal{R}} f dx = 1$. Divide the region between the graph
 $y = f(x)$ and the x -axis into finitely many
 regions, say, S_1, S_2, \dots, S_M , with areas α_1, α_2
 \dots, α_M , respectively. So $\sum_{k=1}^M \alpha_k = \int_{\mathcal{R}} f(x) dx = 1$.
 (See the above figure.)

The composition method to generate a RV $X \sim f$
 is as follows:

step 1 Generate a RV $I \in \{1, \dots, M\}$ with
 the discrete density $(\alpha_1, \alpha_2, \dots, \alpha_M)$.

step 2 Generate (V, W) uniformly in S_I .

step 3 Set $X \leftarrow V$.

If $\{S_k\}$ are mostly regular (e.g., rectangles), then
 step 2 is more efficient.

Generalization: write $f = \sum_{i=1}^M \alpha_i f_i$ as mixture
 of known densities f_i .