

Math 109

August 5, 2016

Question 1 The easiest way to prove an implication of the form $(P \text{ or } Q) \Rightarrow R$ is to prove

- A.** $(P \Rightarrow R)$ and $(Q \Rightarrow R)$.
- B.** its contrapositive $(\text{not } R) \Rightarrow (\text{not } P)$ and $(\text{not } Q)$.
- C.** $[(\text{not } R) \Rightarrow (\text{not } P)]$ and $[(\text{not } R) \Rightarrow (\text{not } Q)]$.
- D.** **B** or **C**, since they're logically equivalent.
- *E.** **A**, **B**, or **C**, since they're all logically equivalent.

Question 2 De Morgan's Law is the logical relationship

A. $\text{not } (P \text{ or } Q) \Leftrightarrow (\text{not } P) \text{ and } (\text{not } Q).$

B. $\text{not } (P \text{ and } Q) \Leftrightarrow (\text{not } P) \text{ or } (\text{not } Q).$

***C.** **A** and **B**: both are part of De Morgan's Law(s).

D. Neither **A** nor **B**: both equivalencies are false.

E. There's no De Morgan's Law! Who was De Morgan anyway?

Question 3 The implication $P \Rightarrow (Q \text{ or } R)$ can be proved by

A. proving $[P \text{ and } (\text{not } Q)] \Rightarrow R$.

B. proving $[P \text{ and } (\text{not } R)] \Rightarrow Q$.

C. cases: that is, by proving $(P \Rightarrow R)$ and $(Q \Rightarrow R)$.

***D.** **A** or **B** since they are logically equivalent.

E. None of the above: $P \Rightarrow (Q \text{ or } R)$ must be proved directly.

Question 4 Theorem: All horses are the same color.

Proof: Clearly, the theorem holds when there is $n = 1$ horse. Suppose the theorem holds for n horses. Given a set of $n + 1$ horses, the inductive hypothesis implies that the 1st through n^{th} and 2nd through $(n + 1)^{\text{st}}$ horses are all the same color. Hence, all $n + 1$ horses are the same color and the theorem follows.

- A.** The theorem is balderdash, since it is easy to exhibit two horses with different colors.
- B.** The proof is faulty because mathematical induction cannot be applied to horses.
- C.** The proof is faulty because because it fails at the inductive step from $n = 1$ to $n = 2$.
- D.** **A** and **B**
- *E.** **A** and **C**

Question 5 Theorem: $n^2 \leq n$ for every positive integer n .

Proof: Clearly, $1^2 \leq 1$. Suppose $k^2 \leq k$ for some positive integer k . If $(k + 1)^2 \leq (k + 1)$, then $k^2 \leq (k + 1)^2 - 1 \leq (k + 1) - 1 = k$. Thus, $(k + 1)^2 \leq k + 1$ must be true.

- A.** The theorem is balderdash. For example, $2^2 > 2$.
- B.** Since the proof is clear, this has become known as the Square Paradox.
- C.** The proof is faulty because it assumes the conclusion as part of the argument.
- D.** **A** and **B**
- *E.** **A** and **C**