

# Cycles in Planar Graphs

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## Abstract

Let  $\mathcal{C}(G)$  denote the set of lengths of cycles in a graph  $G$ , and let  $\mathcal{CP}(n)$  denote the number of distinct subsets  $\mathcal{C}(G) \subset \{1, 2, \dots, n\}$  where  $G$  is a hamiltonian planar graph on  $n$  vertices. In this paper, we prove that

$$\mathcal{CP}(n) \leq 2^{cn}$$

where  $c < 1$  is an absolute positive constant. This compares with the constructive lower bound  $\mathcal{CP}(n) \geq 2^{n/2}$  given by Faudree.

## 1 Introduction

For a graph  $G$ , we write  $\mathcal{C}(G)$  for the set of lengths of cycles in  $G$ , and any subset of  $\{1, 2, \dots, n\}$  which can be realized as  $\mathcal{C}(G)$  for some graph  $G$  on  $n$  vertices is called a cycle set. There has been a great deal of research on the arithmetic properties of cycle sets of graphs relative to their edge-density. Early results include Bondy's Theorem [2], which states that  $\mathcal{C}(G) = \{3, 4, \dots, n\}$  when the minimum degree of  $G$  is greater than  $\frac{n}{2}$  the number of vertices of  $G$ , and a recent theorem of Gould, Haxell and Scott [5], which shows in particular that if  $G$  is a hamiltonian graph whose minimum degree is linear in the number of vertices of  $G$ , then  $\mathcal{C}(G)$  is essentially an interval. In contrast, very little is known for sparse graphs – graphs whose degrees are not linear with respect to the number of vertices in the graph – and there are many open questions which remain. For example, Erdős and Gyárfás asked whether every graph of minimum degree at least three has a cycle of length a power of two, and this remains open. Perhaps one of the most noteworthy questions is whether  $\mathcal{C}(G)$  consists of all even integers in  $\{1, 2, \dots, n\}$  when  $G$  is a hamiltonian  $n$ -vertex graph of minimum degree at least  $c\sqrt{n}$  and  $c$  is a sufficiently large absolute constant. The proof of such a statement seems quite out of reach at present.

In this paper, we consider the number of distinct sets  $\mathcal{C}(G)$  when  $G$  is a hamiltonian planar graph. In [9], it was proved that very few subsets of  $\{1, 2, \dots, n\}$  are the set of cycle lengths of a graph on  $n$  vertices, proving a conjecture of Erdős [4]. We believe that if  $\mathcal{C}(n)$  is the number of distinct cycle sets in  $\{1, 2, \dots, n\}$ , then

$$\lim_{n \rightarrow \infty} \mathcal{C}(n)^{\frac{1}{n}} = \sqrt{2}.$$

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A constructive lower bound of Faudree, consisting of planar graphs, shows that  $C(n) \geq 2^{(n-1)/2}$  for all  $n$ , which supports the limiting behaviour given above. As a first step, we prove the following theorem for planar graphs:

**Theorem 1** *Let  $CP(n)$  denote the number of cycle sets of hamiltonian planar graphs on  $n$  vertices. Then there exists a constant  $c < 1$  such that*

$$CP(n) \leq 2^{cn}.$$

For large  $n$  one can take  $c = \frac{499}{500}$ . It would be interesting to see if  $\lim CP(n)^{1/n}$  differs from  $\lim C(n)^{1/n}$ , as the known constructive lower bounds on  $C(n)$  are all planar. We have not spent much effort on optimizing the value of the constant  $c$  coming from our proofs. It will be quite clear, however, that the techniques we use cannot give the optimal value  $c = \frac{1}{2}$  without some new ideas. In fact, to indicate just how far we are from extending the theorem to all graphs, we raise the following conjecture:

**Conjecture.** *Let  $G$  be a hamiltonian graph of minimum degree at least three. Then there is an absolute positive constant  $a$  such that  $|C(G)| \geq an$ .*

The techniques used to prove Theorem 1 show that the conjecture is true for planar graphs, and one can take  $a = \frac{1}{4}$  in that case. Finally we use the following notation: throughout this paper,  $G$  denotes a finite graph without loops and multiple edges, unless specifically stated otherwise, and  $V(G)$  and  $E(G)$  denote its vertex and edge set respectively. We write  $C(G)$  for the set of lengths of cycles in a graph  $G$ .

## 2 Induced trees and cycles

In this section, we make use duality in planar graphs to prove a variety of statements on the cycle set of a hamiltonian planar graph. Throughout this section,  $G$  is a hamiltonian plane graph with hamiltonian cycle  $C$ . We denote by  $G^*$  the dual graph of  $G$ .

**Lemma 2.1** *Let  $E$  be the set of edges of  $G$  in the interior of  $C$ . Then the set of edges of  $G^*$  which are dual to the edges of  $E$  span a tree in  $G^*$ , and this tree is an induced subtree of  $G^*$  with  $|E| + 1$  vertices.*

**Proof.** It is well-known that if  $D$  is an induced cycle in  $G^*$ , then the set of edges dual to  $E(D)$  form a minimal edge cut in  $G$ . Since no subset of  $E$  is an edge-cut in  $G$ , it follows that the set of edges of  $G^*$  which are dual to the edges of  $E$  span an induced forest in  $G^*$ . The number of edges in this forest is  $|E|$  and the number of vertices in the forest equals the number of faces of  $G$  in the interior of  $C$ , namely  $|E| + 1$ . This implies that the forest is a tree. ■

**Lemma 2.2** *Let  $T$  be any induced subtree of  $G^*$  and suppose  $G^* - V(T)$  is connected. Let  $E_T^*$  be the set of edges of  $G^*$  which are incident with exactly one vertex of  $T$ . Then the set of edges of  $G$  which are dual to the edges of  $E_T^*$  span a cycle of length  $|E_T^*|$  in  $G$ .*

**Proof.** The set  $E_T^*$  is an edge cut in  $G^*$ , and since  $T$  is a tree and  $G^* - V(T)$  is connected,  $E_T^*$  is a minimal edge cut. Therefore the set of edges of  $G$  which are dual to  $E_T^*$  span a cycle of length exactly  $|E_T^*|$  in  $G$ . ■

We now combine the above lemmas to obtain information on the set  $C(G)$ . This is the main lemma of this section. For an induced tree  $U \subset G^*$ , let  $\partial U$  denote the set of edges of  $G^*$  incident with exactly one vertex of  $U$ , and if  $U$  is a subtree of  $T$ , we write  $U < T$ . The sumset of sets  $A$  and  $B$  is  $A + B = \{a + b : a \in A, b \in B\}$ .

**Lemma 2.3** *Let  $G$  be a hamiltonian planar graph with  $m$  edges and  $n$  vertices. Then there exists an induced tree  $T \subset G^*$  with at least  $\frac{1}{2}(m - n) + 1$  vertices such that*

$$\{|\partial U| : U < T\} \subset C(G).$$

*Furthermore, there exist sets  $A$  and  $B$  each of size  $\lceil \frac{1}{4}(m - n) \rceil$  such that  $C(G) \supset A + B$ .*

**Proof.** Let  $C$  be a hamiltonian cycle in  $G$ . Embed  $G$  in the plane so that  $C$  is a convex polygon, and the edges and vertices of  $C$  are the sides and vertices of the polygon, respectively. Then there are  $m - n$  edges of  $G$  which are not in  $C$ , and  $\frac{1}{2}(m - n)$  of these edges lie in the interior or exterior of  $C$ . Via the stereographic projection, the interior and exterior of  $C$  are interchangeable, so we can assume that  $C$  has at least  $\frac{1}{2}(m - n)$  chords in its interior. Let  $E$  be the set of these chords of  $C$ . By Lemma 2.1, the set of edges of  $G^*$  which are dual to the edges of  $E$  span an induced subtree  $T$  of  $G^*$  with at least  $\frac{1}{2}(m - n)$  edges. Then  $T$  has  $\frac{1}{2}(m - n) + 1$  vertices, and since  $G^* - V(U)$  is connected for all  $U < T$ , the cut  $\partial U$  is a minimal edge cut in  $G^*$ . By Lemma 2.2, there is a cycle of length  $|\partial U|$  in  $G$ . This proves the first statement of the lemma. For the second part of the lemma, we partition the edge set of  $T$  into two subsets  $E_1$  and  $E_2$ , each spanning a subtree of  $T$ , and such that  $|E_1| \leq |E_2| \leq |E_1| + 1$ . Let  $T_1 < T$  and  $T_2 < T$  be the subtrees spanned by  $E_1$  and  $E_2$ , and let  $x$  be the unique vertex in common to both  $T_1$  and  $T_2$ . Let  $S_1 < S_2 < \dots < S_{t_1}$  be subtrees of  $T_1$  where  $|V(S_i)| = i$  and  $V(S_1) = \{x\}$  and  $S_{t_1} = T_1$ , and let  $U_1 < U_2 < \dots < U_{t_2}$  be subtrees of  $T_2$ , where  $|V(U_i)| = i$  and  $V(U_1) = \{x\}$  and  $U_{t_2} = T_2$ . Then, with

$$A = \{|\partial S_i| : 1 \leq i \leq t_1\} \quad \text{and} \quad B = \{|\partial U_i| : 1 \leq i \leq t_2\}.$$

we have  $C(G) \supset \{|\partial U| : U < T\} \supset A + B$ . The sets  $A$  and  $B$  have sizes  $t_1$  and  $t_2$ , respectively. By the choice of  $T_1$  and  $T_2$ , we know that  $t_1, t_2 \geq \lceil \frac{1}{4}(m - n) \rceil$ , so  $A$  and  $B$  (or, rather, appropriate subsets of  $A$  and  $B$ ) have the required sizes. This completes the proof. ■

It would be interesting to know how many edges a hamiltonian planar graph  $G$  requires to guarantee  $C(G) = \{3, 4, \dots, n\}$  (such graphs are called pancyclic – see Bondy [2]). It seems possible that every hamiltonian planar graph on  $n$  vertices with more than  $2n$  edges is pancyclic. As a simple application of Lemma 2.3, which is not necessary for the proof of Theorem 1, we obtain a short proof that every hamiltonian maximal planar graph is pancyclic. In fact even more is true – even graphs which are a few edges away from maximal planar are pancyclic – but we omit the more technical details. This result was proved in a different way by Hakimi, Schmeichel and Thomassen [7].

**Corollary 2.4** *Let  $G$  be a hamiltonian maximal planar graph. Then  $G$  is pancyclic.*

**Proof.** A maximal planar graph on  $n$  vertices has  $3n - 6$  edges. Applying Lemma 2.3 with  $m = 3n - 6$ , we see that  $C(G)$  contains  $\{|\partial U| : U < T\}$  for some tree  $T \subset G^*$  with at least  $n - 2$  vertices. If  $U_1 < U_2 < \dots < U_{n-2}$  are subtrees of  $T$  where  $|V(U_i)| = i$ , then

$$|\partial U_1| < |\partial U_2| < \dots < |\partial U_{n-2}|$$

and  $C(G)$  contains  $n - 2$  distinct integers. This implies  $C(G) = \{3, 4, \dots, n\}$ . ■

### 3 Counting sets containing large sumsets

In the last section we showed that  $C(G)$  contains  $A + B$  where  $A$  and  $B$  are large if the number of edges in  $G$  is large. To prove Theorem 1, we will show that very few subsets of  $\{1, 2, \dots, n\}$  contain a set of the form  $A + B$  for some linear-size sets  $A$  and  $B$ . The aim of this section is to quantify this statement, using discrete Fourier analysis. The same method was used in [9]. A recent approach by Green and Ruzsa [6] determines the number of subsets of  $\mathbb{Z}_n$  of the form  $A + A$  when  $n$  is prime. It would be interesting to see if this approach carries over to counting cycle sets, although there are some difficulties to overcome, and one would need a description of the whole cycle set, not only large subsets of it. The following well-known entropy bound on binomial coefficients is needed. The bound follows fairly easily from Stirling's Formula. Let  $a, b \in (0, 1)$ , and let  $H(x) = -x \log x - (1 - x) \log x$ . denote the entropy of  $x \in (0, 1)$ . Then

$$\binom{an}{bn} \leq e^{anH(\frac{b}{a})}. \quad (1)$$

In what follows we will make frequent use of this inequality. We are now ready to prove that there are few sets containing large sumsets:

**Lemma 3.1** *Let  $s(n)$  denote the number of subsets of  $\{1, 2, \dots, n\}$  which contain a set of the form  $A + B$  where  $|A| = an = |B|$ . Then*

$$s(n) \leq n(n - 1)e^{nH(\gamma)} \quad (2)$$

whenever  $\gamma$  satisfies the inequality

$$H(\gamma) \geq \log 2 - 2\gamma^2 a^2. \quad (3)$$

**Proof.** The number of subsets of  $\{1, 2, \dots, n\}$  containing a set of the form  $A + B$  where  $|A| = an = |B|$  is equal to the number of subsets of  $\mathbb{Z}_n$  containing a set of the form  $A + B$ . Let  $S \subset \mathbb{Z}_n$  be a set containing  $A + B$  where  $|A| = an = |B|$  and let  $T = \mathbb{Z}_n \setminus S$ . Let  $\hat{T}$  denote the Fourier transform of the characteristic function  $\chi_T$  of  $T$ , namely

$$\hat{T}(r) = \sum_{t=0}^{n-1} \chi_T(t) e^{\frac{2\pi i r t}{n}}.$$

We claim that if  $|\hat{T}(r)|$  for  $r \neq 0$  achieves a maximum value at  $r_0$ , then

$$|\hat{T}(r_0)| \geq a|T|. \quad (4)$$

To see this, note that  $T \cap (A + B) = \emptyset$ , or equivalently, the equation  $t - x - y = 0$  has no solution with  $t \in T$ ,  $x \in A$  and  $y \in B$ . Since the sum of the roots of unity is zero, one has

$$\sum \chi_T(t) \chi_A(x) \chi_B(y) e^{\frac{2\pi i r(t-x-y)}{n}} = 0$$

where  $\chi_A$  and  $\chi_B$  are the characteristic functions of  $A$  and  $B$  and the sum is over  $r, t, x, y \in \mathbb{Z}_n$ . When  $r = 0$ , the summand above is  $|T||A||B|$ . Therefore taking absolute values and writing  $\hat{A}$  and  $\hat{B}$  for the transforms of  $\chi_A$  and  $\chi_B$ , we obtain

$$\begin{aligned} |T||A||B| &= \left| \sum_{r \neq 0} \chi_T(t) \chi_A(x) \chi_B(y) e^{\frac{2\pi i r(t-x-y)}{n}} \right| \\ &\leq \sum_{r \neq 0} |\hat{T}(r)| |\hat{A}(r)| |\hat{B}(r)| \\ &\leq |\hat{T}(r_0)| \sum_{r=0}^{n-1} |\hat{A}(r)| |\hat{B}(r)| \\ &\leq |\hat{T}(r_0)| \left( \sum_{r=0}^{n-1} |\hat{A}(r)|^2 \right)^{\frac{1}{2}} \left( \sum_{r=0}^{n-1} |\hat{B}(r)|^2 \right)^{\frac{1}{2}} \\ &= |\hat{T}(r_0)| (|A|n)^{\frac{1}{2}} (|B|n)^{\frac{1}{2}}. \end{aligned}$$

In the fourth line we used the Cauchy-Schwarz Inequality, and in the fifth line we used Parseval's Identity, namely

$$\sum_{r=0}^{n-1} |\hat{A}(r)|^2 = n \sum_{r=0}^{n-1} \chi_A(r) = n|A|.$$

Now (4) follows from  $|A| = an = |B|$ . Suppose the elements of  $S$  are chosen uniformly and independently from  $\mathbb{Z}_n$  with probability  $\frac{1}{2}$ . Let  $E$  be the event (4). To finish the proof of the lemma, we have to prove that

$$\mathbb{P}[E] \leq (n-1)n2^{-2 \log_2 e \gamma^2 a^2 n}. \quad (5)$$

To see this, we note that  $|T|$  has a binomial distribution with mean  $\frac{n}{2}$ , and by Hoeffding's Inequality (see Proposition 2.1 and Lemma 2.3 in [9], and McDiarmid [8] for a nice article on concentration inequalities), for  $r \neq 0$ ,

$$\mathbb{P}[|\hat{T}(r)| \geq t] \leq e^{-2t^2/n}.$$

Therefore for any  $r \neq 0$ ,

$$\begin{aligned}
\mathbb{P}[|\hat{T}(r)| \geq a|T|] &= \sum_{t=0}^{n-1} \mathbb{P}[|\hat{T}(r)| \geq ta \cap |T| = t] \\
&\leq \sum_{t=0}^{n-1} \min \left\{ e^{-\frac{2t^2 a^2}{n}}, \binom{n}{t} 2^{-n} \right\} \\
&\leq n \sup_{\gamma \in (0,1)} \min \left\{ e^{-2\gamma^2 a^2 n}, \binom{n}{\gamma n} 2^{-n} \right\} \\
&\leq n \sup_{\gamma \in (0,1)} \min \left\{ e^{-2\gamma^2 a^2 n}, e^{nH(\gamma)} 2^{-n} \right\} \\
&\leq n e^{nH(\gamma)}.
\end{aligned}$$

In the last line we used (3). Since there are  $n - 1$  possible values for  $r_0$ , this gives (5). ▀

## 4 Counting Planar Graphs

To prove Theorem 1 we have to consider the contribution to  $\text{CP}(n)$  of those cycle sets of hamiltonian planar graphs which have few edges (this is quantified later on). We will prove that the number of hamiltonian planar graphs with  $m$  edges and  $n$  vertices is much less than  $2^n$  whenever  $m$  is not too much larger than  $m$ , and this clearly gives an upper bound on the contribution of these graphs to  $\text{CP}(n)$ . The next lemma counts the number of hamiltonian planar graphs with  $m$  edges. The problem of counting planar blocks was solved by Bender, Gao and Wormald [1] using analytic techniques. The bound we give below is weaker, however the proof is very much simpler and suffices for our purposes.

**Lemma 4.1** *Let  $u(m, n)$  denote the number of unlabelled hamiltonian planar graphs with  $m$  edges and  $n$  vertices, where  $m > n$ . Then*

$$u(m, n) \leq (m - n) e^{2mH(1 - \frac{n}{m})} 2^{2m - 2n} \quad (6)$$

**Proof.** Let  $G$  be a hamiltonian planar graph. Draw a hamiltonian cycle  $C$  of  $G$  in the plane as a convex polygon, and all remaining edges of  $G$  as chords of  $C$ . Then the number of graphs

$G$  is at most the number of ways to place chords in the interior of  $C$  times the number of ways to place chords in the exterior of  $C$ . Let  $G_0$  denote the graph obtained by deleting all exterior chords of  $C$ . Then we may contract pairs of adjacent vertices of  $C$  until we obtain a triangulation  $H_0$  with the same number of edges as  $G_0$ . Let  $i$  denote the number of chords of  $C$  in  $G_0$ , and let  $C_0$  denote the hamiltonian cycle of  $H_0$  corresponding to  $C$ . Then  $C_0$  has length  $i + 3$ , and we may order the vertices of  $C_0$  as  $(v_1, v_2, \dots, v_{i+3})$  in the clockwise orientation of  $C_0$ . We will show that  $H_0$  is uniquely determined by a pair  $(H_0, \sigma)$ , where  $\sigma$  is a composition of  $n$  into  $2i$  parts. For a given triangulation  $H_0$  with hamiltonian cycle  $C_0$ , the composition  $\sigma$  is defined as follows. First let  $F$  denote the graph obtained from  $H_0$  by deleting all edges  $\{v_h, v_{h+1}\} \in E(C_0)$  such that  $v_h$  has degree two in  $H_0$ . For each  $\{v_h, v_{h+1}\} \in E(C_0) \cap E(F)$ , assign a non-negative integer  $x_h$ , denoting the number of vertices inbetween  $v_h$  and  $v_{h+1}$  in  $C$ . Define

$$\tau = (x_h : \{v_h, v_{h+1}\} \in E(C_0) \cap E(F)).$$

Now we list edges incident with  $v_h$  counterclockwise around  $v$ , say  $e_{1,h}, e_{2,h}, \dots, e_{d_h,h}$ . Let  $y_{g,h}$  denote the number of vertices in the subpath of  $C_0$  joining the endpoints of  $e_{g,h}$  and  $e_{g+1,h}$  which were contracted into  $v_h$ . Then we obtain a sequence  $\sigma_h = (y_{1,h}, y_{2,h}, \dots, y_{d_h,h})$  for each  $h \geq 1$ . Finally,  $\sigma$  is defined by

$$\sigma = \tau\sigma_1\sigma_2 \dots \sigma_{i+3}$$

and  $\sigma$  has length  $2i$ . This is illustrated in Figure 1, with  $i = 6$  and  $\sigma = (0, 2, 2, 1, 2, 2, 0, 1, 0, 0, 2, 1)$ . The dotted lines denote edges of  $C_0$  which are not in  $F$ .

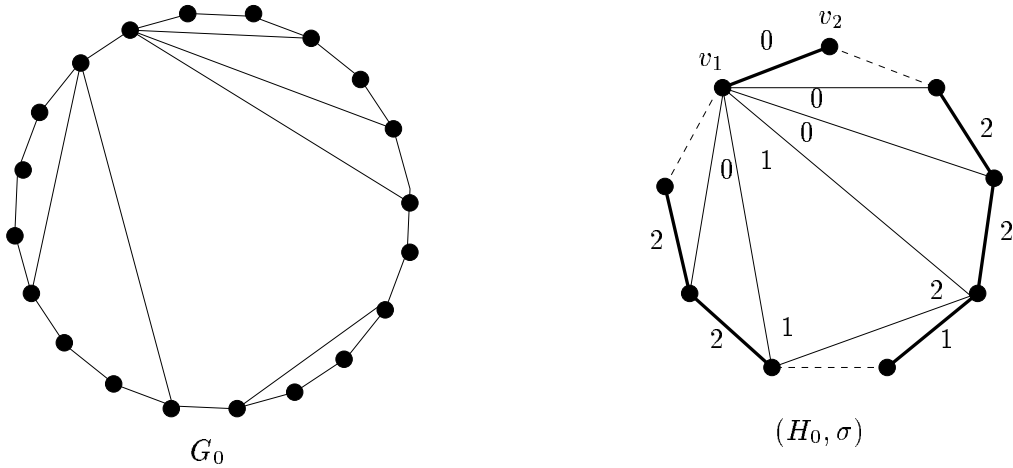


Figure 1.

Now we count the number of pairs  $(H_0, \sigma)$ . The number of compositions of  $n$  into  $2i$  non-negative parts is exactly  $\binom{n+2i-1}{2i-1}$ . Euler determined the number of triangulations of an  $r$ -gon, and that number is exactly the Catalan number  $\frac{1}{r-1} \binom{2r-4}{r-2}$ . Since  $C_0$  is an  $(i + 3)$ -gon, the number  $f(i)$  of graphs  $G_0$  with  $i$  edges in the interior of  $C$ , using (1), is at most

$$\binom{n+2i-1}{2i-1} \cdot \frac{1}{i+2} \binom{2i+2}{i+1} < \exp \left\{ (n+2i)H\left(\frac{2i}{n+2i}\right) \right\} \cdot 2^{2i}.$$

The same approach gives an upper bound on the number of graphs  $G_1$  consisting of  $C$  and all exterior chords of  $C$ . Furthermore, the number of planar graphs with  $i$  edges in the interior of  $C$  and  $j = m - n - i$  edges in the exterior of  $C$  is at most the product of  $f(i)$  and  $f(j)$ . This implies that  $u(m, n)$  is at most the sum over all  $i$  and  $j$  adding up to  $n$  of  $f(i)f(j)$ . It follows that

$$u(m, n) < \sum_{\substack{i+j \\ =m-n}} \exp \left\{ (n+2i)H\left(\frac{2i}{n+2i}\right) + (n+2j)H\left(\frac{2j}{n+2j}\right) \right\} \cdot 2^{2i+2j}.$$

The summand is maximized when  $2i = 2j = m - n$ , by symmetry, and it follows that  $u(m, n)$  is at most  $m - n$  times that value, giving (6).  $\blacksquare$

## 5 Proof of Theorem 1

Let  $\mathcal{P}$  and  $\mathcal{Q}$  denote the set of all hamiltonian planar graphs on  $n$  vertices with at most  $(1 + \alpha)n$  edges and more than  $(1 + \alpha)n$  edges, respectively, where  $\alpha > 0$  is to be chosen later. According to the bound (6) on the number of graphs in  $\mathcal{P}$  in Lemma 4.1, we have

$$|\{C(G) : G \in \mathcal{P}\}| \leq |\mathcal{P}| \leq e^{2(1+\alpha)n \cdot H(1-\frac{1}{1+\alpha})} 2^{2\alpha n}. \quad (7)$$

Next we consider  $\mathcal{Q}$ . For each planar graph  $G \in \mathcal{Q}$ , Lemma 2.3 shows that  $C(G)$  contains  $A + B$  for some sets  $A$  and  $B$  in  $\{1, 2, \dots, n\}$  where  $|A| = |B| = \lceil \frac{\alpha n}{4} \rceil$  each. By Lemma 3.1,

$$|\{C(G) : G \in \mathcal{Q}\}| \leq s(n) \leq n(n-1)e^{nH(\gamma)} \quad (8)$$

where  $\gamma$  is defined by (3) with  $a = \frac{\alpha}{4}$ . Adding (7) and (8) gives an upper bound on  $\text{CP}(n)$ . Finally we determine the value of  $\alpha$  for which the logarithms of (7) and (8) are asymptotically equal as  $n$  tends to infinity. This is when

$$\frac{1}{4}\alpha^2\gamma^2 = 2\alpha \log \frac{\alpha}{\alpha+1} + 2 \log \frac{1}{\alpha+1} - 2\alpha(\log 2) + \log 2.$$

Then we solve for the value of  $\alpha$  using (3), to obtain  $0.081779 < \alpha < 0.08178$ . In particular

$$\text{CP}(n) < e^{0.6913n} < 2^{499n/500}$$

provided  $n$  is sufficiently large. This completes the proof of Theorem 1.  $\blacksquare$

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