

## On Partitions of a Finite Set

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### ABSTRACT

A pair of partitions  $\pi_1, \pi_2$ , of a finite set  $S$  into disjoint non-empty subsets will be called *conjugate* if for each  $s \in S$ , the ordered pair  $(v_1(s), v_2(s))$  determines  $s$ , where  $v_i(s)$  denotes the cardinality of the subset of  $\pi_i$  to which  $s$  belongs. In this note we show that  $S$  has a pair of conjugate partitions if and only if the cardinality of  $S$  is not equal to 2, 5, or 9. Partitions of this type provide a short solution to a problem arising in circuit theory.

### INTRODUCTION

Suppose we have a cable consisting of  $n$  indistinguishable wires with terminals at two points  $A$  and  $B$ , and suppose for each terminal at  $A$  it is desired to identify its mate at  $B$ . We shall assume that the only operations available for making such an identification are interconnecting sets of terminals at one end and testing for current flow in the terminals at the other end. For example, if all terminals at  $A$  are connected together, then a current can flow between any two terminals at  $B$ . Without this assumption, the desired identification would present no problem, since if we denote the terminals at  $A$  by  $A_i$ ,  $1 \leq i \leq n$  (and similarly for  $B$ ), then we simply test to see if a current can flow between  $A_1$  and  $B_1$ ,  $A_1$  and  $B_2, \dots$ , until we find a  $B_{i_1}$  such that a current can flow between  $A_1$  and  $B_{i_1}$  and consequently we know  $A_1$  and  $B_{i_1}$  represent the same wire. We then use the same procedure on  $A_2$ , etc. For long cables, we shall restrict ourselves further to procedures of the following type:

Certain connections are made at  $A$ . We then go to  $B$  and make tests and, using the test results, certain connections. We finally come back to  $A$ , disconnect the connections initially made, and perform further tests. The information now in hand should be enough to determine for each  $j$  the terminal pairs  $A_j$  and  $B_{i_j}$  of wire  $j$ .

The following ingenious algorithm for solving this problem is due to my colleague K. C. Knowlton. Before presenting the general solution, a typical specific example will be given. Consider the case  $n = 6$ . Define the partitions  $P: \{1, 2, 3\}, \{4, 5\}, \{6\}$ , and  $P': \{1, 4, 6\}, \{2, 5\}, \{3\}$ , of the integers  $\{1, 2, 3, 4, 5, 6\}$ . If we let  $f(i)$  denote the number of elements in the set of  $P$  which contains  $i$  and  $f'(i)$  the similar function for the partition  $P'$ , then we have the table

$i$	$(f(i), f'(i))$
1	(3,3)
2	(3,2)
3	(3,1)
4	(2,3)
5	(2,2)
6	(1,3)

We note that in the table all the ordered pairs  $(f(i), f'(i))$  are *distinct*. To identify the terminals in our cable of six wires we first label those at  $A$  by  $A_1, A_2, \dots, A_6$  and connect them as shown in Figure 1(a). We now test at  $B$  to decide which wires have been connected at  $A$ .

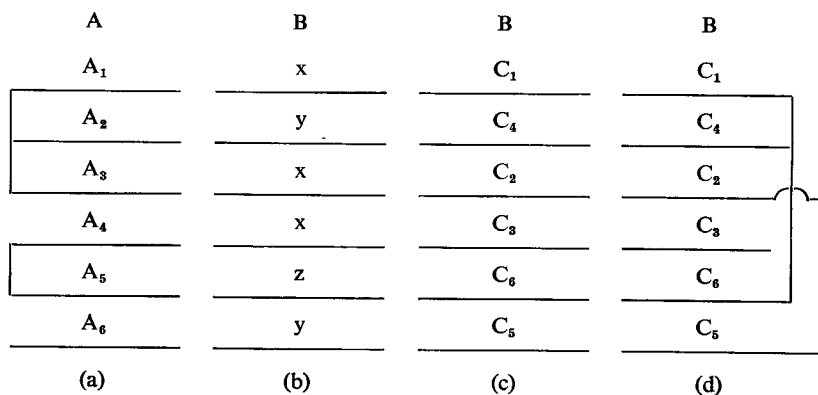


FIGURE 1

Suppose, for example, we find that we have the situation indicated in Figure 1(b) where two wires with the same symbol are joined at  $A$ . We label the  $x$ 's with  $C_1, C_2$  and  $C_3$  (arbitrarily), the  $y$ 's with  $C_4$  and  $C_5$  and the  $z$  with  $C_6$ , say, as in Figure 1(c). We next connect the terminals at  $B$  together according to the partition  $P'$ , i.e., as in Figure 1(d). Finally we go back to  $A$ , disconnect the connections initially made there, and test to decide how many wires a given wire at  $A$  has been joined to at  $B$ . For example, suppose we find that wire  $A_1$  is now connected to exactly one other wire at  $B$ . Since we know that initially  $A_1$  was in a set of *three* wires which were connected at  $A$  and now it is in a set of *two* wires which are connected at  $B$  and, since the pair of  $(f(i), f'(i)) = (3, 2)$  occurs only for  $i = 2$ , then we can conclude that  $A_1$  must correspond to  $C_2$ . Similarly if it happens that  $A_5$  is now connected to two other wires at  $B$  (so that  $A_5$  belongs to a set of *three* wires which have been joined) then  $A_5$  must correspond to  $C_4$ , etc.

The general solution may be described in the following way. Let  $I_n$  denote the set of integers  $\{1, 2, \dots, n\}$ . Suppose there exists a pair of partitions of  $I_n$ , say  $P: P_1, P_2, \dots, P_k$  and  $P': P'_1, P'_2, \dots, P'_{k'}$ , such that, if  $f(j)$  denotes the cardinality of the subsets  $P_i$  which contains  $j$  (with  $f'(j)$  defined similarly), then the map  $j \rightarrow (f(j), f'(j))$  is a 1-1 mapping of  $I_n$  into  $I_n \times I_n$ . We shall call such a pair of partitions *conjugate*. For  $1 \leq j \leq k$ , we first connect all wires together at  $A$  which have subscripts that belong to the same  $P_j$ . We then go to  $B$  and, by suitable testing, determine the subsets  $S_1, \dots, S_k$  of the  $B_i$  which have been connected together at  $A$ . We next relabel the  $B_i$  by  $C_1, C_2, \dots, C_n$  so that for any  $S_r$  the set of indices of the  $C_i$  which occur in  $S_r$  is exactly one of the  $P_j$ . Now we connect the  $C_i$  together to form subsets  $T_j$ ,  $1 \leq j \leq k'$ , in such a way that, for any  $j$ , the set of indices of the  $C_i$  which are in  $T_j$  is just  $P'_j$ . Finally we go back to  $A$ , disconnect all connections previously made there, and by suitable testing decide which  $A_i$  have been connected at  $B$ . We are now in a position to determine which labels represent the same wire. For if we take any wire, say  $A_u$ , we know that at  $B$  it belongs to a  $T_j$  which has, say,  $p$  elements (where we can determine  $p$ ). Since we also know the cardinality of the  $S_i$  to which  $A_u$  belongs, say  $q$ , then, by the way the  $S_i$  and  $T_i$  were constructed and by the hypothesis that all the pairs  $(p, q) = (f(u), f'(u))$  are distinct, we can determine the unique  $C_{i_u}$  and hence the  $B_{j_u}$  such that  $A_u$  and  $B_{j_u}$  represent the same wire.

It is the purpose of this note to prove that pairs of conjugate parti-

tions exist for  $I_n$  if and only if  $n \neq 2, 5, \text{ or } 9$ . We also give a simple construction of a pair of conjugate partitions for  $I_n$  for each  $n \neq 2, 5, \text{ or } 9$ .

### SOME NECESSARY CONDITIONS

As usual let  $|A|$  denote the cardinality of the set  $A$ . If  $P: P_1, P_2, \dots, P_k$  is a partition of  $I_n$  let  $C(P)$  denote the set  $\{|P_1|, |P_2|, \dots, |P_k|\}$  where we shall assume from now on that  $|P_1| \leq |P_2| \leq \dots \leq |P_k|$ .

LEMMA 1. *If  $P$  and  $P'$  are conjugate partitions of  $I_n$  then  $|P_k| = |P'_k|$ .*

PROOF. Suppose  $|P_k| = m$ . Since  $P$  and  $P'$  are conjugate then  $C(P')$  must contain at least  $m$  distinct elements. Consequently  $|P'_k| \geq m$  (since we assume that  $|P'_1| \leq \dots \leq |P'_k|$ ). Applying the same argument to  $P'$  we see that  $|P_k| \leq |P'_k| \leq |P_k|$  and the lemma follows.

If  $P$  is a partition of  $I_n$  for which there exists a partition  $P'$  of  $I_n$  such that  $P$  and  $P'$  are conjugate then we shall call  $P$  *admissible*.

LEMMA 2. *If  $P$  is admissible and  $|P_k| = m$  then  $C(P) = \{1, 2, \dots, m\}$ .*

PROOF. Suppose there exists  $j$  such that  $1 \leq j \leq m$  and  $j \notin C(P)$ . Since  $|P'_k| = m$  by Lemma 1 then we must have  $|C(P)| \geq m$  (since  $P$  is admissible). Therefore  $|P_k| \geq m + 1$ , which is a contradiction.

LEMMA 3. *Suppose  $P$  is admissible,  $|P_k| = m$  and  $n_j$  denotes the number of  $P_i$  such that  $|P_i| = j$ . Then  $n_j \leq [m/j]$  (where  $[x]$  denotes the greatest integer not exceeding  $x$ ).*

PROOF. Suppose there exists  $j$  such that  $n_j > [m/j]$  and let  $P'$  be a partition of  $I_n$  which is conjugate to  $P$ . Since  $r \in P_a, s \in P_b$  and  $|P_a| = |P_b| = j$  imply that  $f'(r) \neq f'(s)$  then we must have  $|C(P')| \geq j \cdot n_j$ . But by Lemmas 1 and 2 (since  $P'$  is also admissible)

$$m = |C(P)| = |C(P')| \geq j \cdot n_j \geq j([m/j] + 1) > j \cdot m/j = m$$

which is a contradiction. This proves the lemma.

We combine these lemmas to obtain

THEOREM 1. *If  $P$  is an admissible partition of  $I_n$  and  $|P_k| = m$  then we have*

$$\frac{m(m+1)}{2} \leq n \leq \sum_{j=1}^m j \left\lfloor \frac{m}{j} \right\rfloor. \quad (1)$$





and project two  $x$ 's and add the extra  $x$  (since  $r = 3 = m - 1$ ) to obtain

	4	3	2	1	1	1	1
4	$x$	$x$	$x$	$x$			
3	$x$	$x$	*		$x$		
2	$x$	*				$x$	
1	$x$						
1			$x$				
1		$x$					
1							$x$

We arbitrarily replace the  $x$ 's by the elements of  $I_{13}$  to form

	4	3	2	1	1	1	1
4	1	2	3	4			
3	5	6	*		7		
2	8	*				9	
1	10						
1			11				
1		12					
1							13

from which we generate the conjugate partitions of  $I_{13}$ :

$$P : \{4\}, \{7\}, \{9\}, \{13\}, \{3, 11\}, \{2, 6, 12\}, \{1, 5, 8, 10\}$$

$$P' : \{10\}, \{11\}, \{12\}, \{13\}, \{8, 9\}, \{5, 6, 7\}, \{1, 2, 3, 4\}.$$

The only  $n$  for which  $I_n$  has not been shown to have conjugate partitions are those of the form  $\Delta(m) - 1$  and, indeed, we have already noted that no such partitions exist for  $I_2, I_5,$  or  $I_9$ . We fill this gap with

**THEOREM 3.** *If  $n = \Delta(m) - 1$  for  $m \geq 4$  then there exist conjugate partitions for  $I_n$ .*

PROOF. We start with the array

	$m - 1$	$m - 2$	$\cdots$	$3$	$2$	$\overbrace{\hspace{2cm}}^{m - 2}$				
						1	1	$\cdots$	1	2
$m - 1$	$x$	$x$	$\cdots$	$x$	$x$	$x$				
$m - 2$	$x$	$x$		$x$	$x$					
$\vdots$	$\vdots$			$x$						
$3$	$x$	$x$								
$2$	$x$	$x$								
$\left\{ \begin{array}{l} 1 \\ 1 \\ \vdots \\ 1 \\ 2 \end{array} \right.$	$x$									

and project the  $x$ 's which have coordinates  $(j, m - j)$  for  $2 \leq j \leq m - 2$ . Next we project the  $x$  at  $(3, m - 2)$  and place an additional  $x$  at the intersection of the last row and column (so that it has coordinates  $(2, 2)$ ). Note that the new points formed from the projection of the  $x$  at  $(3, m - 2)$  have coordinates  $(2, m - 2)$  and  $(3, 2)$  which are distinct from the coordinates of any other  $x$ 's in the array (since  $m \geq 4$  and the  $x$  which was originally at  $(2, m - 2)$  has been projected). It is easily checked that all  $x$ 's in the array have distinct coordinates so that by replacing the  $x$ 's by the elements of  $I_n$  we can form conjugate partitions of  $I_n$  and the theorem is proved.

### CONCLUDING REMARKS

It is interesting to note that, by extending the construction used in Theorem 3, it is possible to form an admissible partition  $P$  of  $I_n$  with  $|P_k| = m$  and  $n = J(m)$  thus achieving the upper bound derived in Theorem 1. However, it is not clear that if  $n$  is any integer such that  $\Delta(m) \leq n \leq J(m)$  then there exists an admissible partition of  $I_n$  with  $|P_k| = m$ . Since Lemmas 2 and 3 show that for any admissible partition  $P$  with  $|P_k| = m$ , we must have  $1 \leq n_j \leq [m/j]$  then it might be conjectured that any partition  $P$  with  $1 \leq n_j \leq [m/j]$  is admissible. This is not the case, however, as the following example shows. Let



$n = 28$  and choose  $P$  so that  $n_1 = 1$ ,  $n_2 = 3$ ,  $n_3 = 2$ ,  $n_4 = 1$ ,  $n_5 = 1$  and  $n_6 = 1$ . Suppose  $P'$  is a partition conjugate to  $P$  and let  $F$  denote  $\{(f(j), f'(j)): 1 \leq j \leq 28\}$ . We shall derive a contradiction. Since the pair  $(1,6)$  must belong to  $F$  then  $n_2' \leq 2$  and  $n_3' = 1$ . But  $(6,3)$ ,  $(3,3)$ , and  $(2,3) \in F$  so that neither  $(5,3)$  nor  $(4,3)$  belong to  $F$ . Hence  $(5,1)$ ,  $(5,2)$ ,  $(5,4)$ ,  $(5,5)$ , and  $(5,6) \in F$  and therefore, since  $(6,4)$ ,  $(5,4)$ ,  $(3,4)$ , and  $(2,4) \in F$  then  $(4,4) \notin F$ . But we must have  $(6,2)$ ,  $(5,2)$ ,  $(3,2)$ , and  $(2,2) \in F$  so that  $(4,2) \notin F$ . Thus we have shown that  $(4,4)$ ,  $(4,3)$ , and  $(4,2)$  cannot belong to  $F$  which is impossible. Hence  $P$  is not an admissible partition of  $I_{28}$ .

It may be shown that, by using this procedure, the maximum number of circuit tests which must be made for the identification is essentially  $2n \log_2 n$ . On the other hand, information theory arguments show that the number of tests must be at least  $\log_2(n!) \approx n \log_2 n$ .

It should be remarked that the function

$$J(m) = \sum_{j=1}^m j \left[ \frac{m}{j} \right]$$

can easily be put into the equivalent form

$$J(m) = \sum_{j=1}^m \sum_{d|j} d.$$

In this form  $J(m)$  is recognized as a well-studied number theoretic function about which statements such as

$$J(m) \sim \frac{\pi^2}{12} m^2 + O(m \log m)$$

can be made (cf. [1]).

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#### REFERENCE

1. W. J. LEVEQUE, *Topics in Number Theory*, Addison-Wesley, Reading, Mass., 1961, p. 121.