

On hypergraphs having evenly distributed subhypergraphs

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Abstract

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1. Introduction and preliminaries

Let k be a fixed positive integer. By a k -uniform hypergraph G , or k -graph for short, we mean a pair (V, E) , where $V = V(G)$ is a set, called the *vertices* of G , and $E = E(G)$ is a subset of $\binom{V}{k}$, the k -element subsets of V , called the *edges* of G (for a full discussion of hypergraphs, see [1]). If V has cardinality $|V| = n$, we denote this by writing $G = G(n)$.

For a k -graph $G' = (V', E')$, we say that G' is an *induced subgraph* of G , written as $G' < G$, if there is a mapping $\lambda: V' \rightarrow V$ such that $X \in E'$ if and only if $\lambda(X) \in E$ (where for $X \in \binom{V'}{k}$, $\lambda(X) := \bigcup_{x \in X} \lambda(x)$). We denote by $\#\{G' < G\}$ the number of such (ordered) mappings.

If $\mathcal{G} = \{G(n) | n \rightarrow \infty\}$ is a family of k -graphs, we say that \mathcal{G} *satisfies* $U(r)$ if, for each k -graph $G'(r)$ on r vertices,

$$\#\{G'(r) < G(n)\} = (1 + o(1))n^r / 2^{\binom{r}{k}}, \quad n \rightarrow \infty. \quad (1)$$

Thus, \mathcal{G} satisfies $U(r)$ if and only if all r -vertex k -graphs occur as (ordered) induced subgraphs of $G(n)$ asymptotically equally often as $n \rightarrow \infty$.

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In particular, if each $G(n)$ is a *random k -graph* $G_{1/2}(n)=(V(n), E(n))$ on n vertices, i.e., each $X \in \binom{V(n)}{k}$ is chosen as an edge of $G_{1/2}(n)$ independently with probability $1/2$, then the corresponding family $\mathcal{G}_{1/2}$ almost certainly satisfies $U(r)$ for any fixed r (i.e., satisfies $U(r)$ with probability tending to 1 as $n \rightarrow \infty$).

It is not difficult to see that if \mathcal{G} satisfies $U(r)$ then \mathcal{G} also satisfies $U(s)$ for any $s \leq r$. On the other hand, it is perhaps unexpected that it is possible to reverse this implication once s is as large as $2k$. More precisely, it was shown in [2] that:

$$\text{If } \mathcal{G} \text{ satisfies } U(2k) \text{ then } \mathcal{G} \text{ satisfies } U(r) \text{ for any fixed } r. \tag{2}$$

Families \mathcal{G} satisfying (2) have been termed *quasi-random*, since it is known that they must necessarily also satisfy a large collection of other properties all shared by families of random k -graphs (for details, see [2-4]).

However, it was noted in [2] that (2) is no longer valid if $U(2k)$ is replaced by $U(k+1)$. The main purpose of this note is to close this gap completely, by showing that (2) no longer holds even if we assume \mathcal{G} satisfies $U(2k-1)$. More generally, for each s , with $k \leq s \leq 2k-1$, there are families \mathcal{G}_s which satisfy $U(s)$ but not $U(s+1)$. A less direct proof for this construction appears in [3].

2. The main construction

If $G=(V, E)$ is a k -graph, we let

$$\chi = \chi_G: \binom{V}{k} \rightarrow \{0, 1\} \text{ be the edge function for } G,$$

defined for $X \in \binom{V}{k}$ by

$$\chi(X) = \begin{cases} 1 & \text{if } X \in E, \\ 0 & \text{if } X \notin E. \end{cases}$$

For $a \geq 0$, we define the *coboundary operator* $\delta^{(a)}$ mapping k -graphs on V to $(k+a)$ -graphs on V as follows. If $G=(V, E)$ is a k -graph with edge function χ , then $\delta^{(a)}G=(V, E^{(a)})$ is a $(k+a)$ -graph with edge function $\chi^{(a)}$, given, for $Y \in \binom{V}{k+a}$, by

$$\chi^{(a)}(Y) \equiv \sum_{X \in \binom{Y}{k}} \chi(X) \pmod{2}. \tag{3}$$

Thus, Y is an edge of $\delta^{(a)}G$ if and only if Y contains an *odd* number of edges X of G as subsets.

For $1 \leq j \leq k-1$, choose a random j -graph $G_{1/2}^{(j)}$ on V and ‘lift’ it to a k -graph $G_j := \delta^{(k-j)}G_{1/2}^{(j)}$ on V with edge function χ_j . Next, form the ‘symmetric difference’ k -graph $G^*(n)=(V, E^*(n)) = \bigvee_{j=1}^{k-1} G_j$ with edge function χ^* , defined by

$$\chi^*(X) \equiv \sum_{j=1}^{k-1} \chi_j(X) \pmod{2} \quad \text{for } X \in \binom{V}{k}.$$

Theorem 2.1. For almost all choices of $G_{1/2}^{(j)}$, $\mathcal{G}^* = \{G^*(n) | n \rightarrow \infty\}$ satisfies $U(2k-1)$ but not $U(2k)$, provided $k \neq 2^t$, $t \geq 0$.

Proof. Consider an arbitrary fixed set $W = \{w_1, w_2, \dots, w_{2k-1}\}$ of $2k-1$ vertices of V . Form a matrix M with rows indexed by $X \in \binom{W}{k}$, and columns indexed by $Y_j \in \binom{W}{j}$, $1 \leq j \leq k-1$, with the (X, Y_j) -entry $M(X, Y_j)$ defined to be 1 if $Y_j \subset X$, and 0 otherwise. We can regard each column $C(Y_j)$ of M as a function mapping $\binom{W}{k}$ to $\{0, 1\}$ by defining

$$C(Y_j)(X) = M(X, Y_j), \quad X \in \binom{W}{k}.$$

It is easy to see that

$$\chi_j \equiv \sum_{Y_j \in \binom{W}{j}} C(Y_j) \pmod{2}.$$

so that

$$\chi^* \equiv \sum_{j=1}^{k-1} \sum_{Y_j \in \binom{W}{j}} C(Y_j) \pmod{2}.$$

We now apply a result of Wilson [6] (see also [5]) which asserts that, when $k \neq 2^t$, M has mod 2 rank equal to $\binom{2k-1}{k}$. Actually, Wilson's result implies that if we adjoin a column of all 1's to form an augmented matrix M^+ , then for any prime p , M^+ has mod p rank equal to $\binom{2k-1}{k}$. However, when $k \neq 2^t$, then some i , with $1 \leq i \leq k-1$, has $\binom{k}{i}$ odd. Summing all the columns $C(Y_i)$, $Y_i \in \binom{V}{i}$, gives us a column of all 1's (mod 2), from which it follows that M itself has mod 2 rank equal to $\binom{2k-1}{k}$.

Now, as W ranges over all $(2k-1)$ -element subsets of V , since the edges of the various corresponding $G_{1/2}^{(j)}$ are chosen independently and uniformly, then an easy argument shows that almost certainly each of the possible $\binom{2k-1}{k}$ $(0, 1)$ -vectors occurs $(1 + o(1))n^{2k-1} / 2^{\binom{2k-1}{k}}$ times as $n \rightarrow \infty$. But this just means that for almost all choices of the $G_{1/2}^{(j)}$, each of the possible k -graphs $G(2k-1)$ on $2k-1$ vertices occurs $(1 + o(1))n^{2k-1} / 2^{\binom{2k-1}{k}}$ times as an induced subgraph of $G^*(n)$ as $n \rightarrow \infty$. This implies that \mathcal{G}^* satisfies $U(2k-1)$, as claimed.

To show that \mathcal{G}^* does not satisfy $U(2k)$, we do the following. Let $Z = \{z_1(0), z_1(1), \dots, z_k(0), z_k(1)\}$ be an arbitrary $2k$ -element subset of V . Consider the inclusion matrix \bar{M} with rows indexed by $X \in \binom{Z}{k}$, and columns indexed by $Y_j \in \binom{Z}{j}$, $1 \leq j \leq k-1$. Let us restrict our attention to the 2^k rows indexed by all the k -sets of the form $Z(\varepsilon_1, \dots, \varepsilon_k) = \{z_1(\varepsilon_1), \dots, z_k(\varepsilon_k)\}$, where $\varepsilon_i \in \{0, 1\}$, $1 \leq i \leq k$. For a fixed $Y_j \in \binom{Z}{j}$, since $j \leq k-1$, there must exist at least one index i such that $z_i(0) \notin Y_j$, $z_i(1) \notin Y_j$. Thus, $Z(\varepsilon_1, \dots, \varepsilon_{i-1}, 0, \varepsilon_{i+1}, \dots, \varepsilon_k) \supset Y_j$ if and only if $Z(\varepsilon_1, \dots, \varepsilon_{i-1}, 1, \varepsilon_{i+1}, \dots, \varepsilon_k) \supset Y_j$. This implies that for any column of \bar{M} , the total number of 1's in the 2^k rows indexed by the $Z(\varepsilon_1, \dots, \varepsilon_k)$ is always even. Hence, this also holds for the mod 2 sum of any set of columns of \bar{M} . Consequently, \mathcal{G}^* contains no indexed subgraph on a $2k$ -set Z in which an odd number of k -sets $Z(\varepsilon_1, \dots, \varepsilon_k)$ are edges. This shows that \mathcal{G}^* does not satisfy $U(2k)$, and Theorem 2.1 is proved. \square

In the case that $k=2^t$, an additional step is required. As before, we first construct the k -graphs $G^*(n)=(V_n, E)$. We then take the complement $\bar{G}^*(n)=(V'_n, E')$ on a disjoint vertex set V'_n , and form the k -graph $\hat{G}(2n):=(V_n \cup V'_n, \hat{E})$ by defining \hat{E} to be $E \cup E'$, together with a random selection of all the k -sets X which intersect both E and E' . That is, each such X is chosen independently with probability $1/2$ to be an edge of $\hat{G}(2n)$.

Theorem 2.2. *For almost all choices of $G_{1/2}^{(j)}$, $\mathcal{G} = \{\hat{G}(2n) | n \rightarrow \infty\}$ satisfies $U(2k-1)$ but not $U(2k)$.*

Proof. The case not covered by Theorem 2.1 is when $k=2^t$, which we now assume. By the previously mentioned result of Wilson, if $G(2k-1)$ is a k -graph on $2k-1$ vertices then

$$\begin{aligned} & \# \{G(2k-1) < G^*(n)\} \\ &= \begin{cases} 2(1+o(1))n^{2k-1}/2^{\binom{2k-1}{k}} & \text{if } G(2k-1) \text{ has an even number of edges,} \\ 0 & \text{if } G(2k-1) \text{ has an odd number of edges.} \end{cases} \end{aligned}$$

Since $\binom{2k-1}{k}$ is odd for $k=2^t$, the situation is reversed for the complement $\bar{G}^*(n)$. This implies that \mathcal{G}^* satisfies $U(2k-1)$.

To see that \mathcal{G}^* does not satisfy $U(2k)$, consider the k -graph $H = H((2k-1)^2)$ formed from disjoint copies of $H_i(2k-1) = (W_i, E_i)$, $1 \leq i \leq 2k-1$, where each $H_i(2k-1)$ is a complete k -graph on $2k-1$ vertices, i.e., $|W_i| = 2k-1$ and $E = \binom{W_i}{k}$. We claim:

$$\# \{H < \hat{G}(2n)\} = 0 \quad \text{for all } n. \quad (4)$$

To see this, suppose the contrary. Note that we must have $W_i \not\subseteq V_n$ for $1 \leq i \leq 2k-1$, since otherwise $H_i(2k-1) < G^*(n)$, which is impossible, because each $H_i(2k-1)$ has an odd number $\binom{2k-1}{k}$ of edges. Thus, for each i there is some $w_i \in V'_n$, $1 \leq i \leq 2k-1$. However, the k -graph induced by the vertex set $\{w_1, \dots, w_{2k-1}\}$ has no edges, which contradicts the fact that $\bar{G}^*(n)$ only has induced subgraphs on $2k-1$ vertices having an odd number of edges. This proves (4). Finally, by (2) this implies that \mathcal{G} does not satisfy $U(2k)$, and the proof is complete. \square

We remark that essentially the same arguments can be applied for any s with $k \leq s \leq 2k-1$, i.e., showing that families \mathcal{G}_s exist satisfy $U(s)$ but not $U(s+1)$. We omit the details.

3. Concluding remarks

It would be interesting to know whether, in fact, the cases $k=2^t$ are inherently different, or whether this is simply an artifact of the approach we have taken. For

example, we do not know, for any $k=2^t$, whether there exists a family $\mathcal{G} = \{G(n) | n \rightarrow \infty\}$ of k -graphs satisfying $U(2k-1)$ but for which, for some $H(2k)$,

$$\#\{H(2k) < G(n)\} = 0 \quad \text{for all } n.$$

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