

HW1 Solution of Math170A

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1.1.25

Proof. Let $C^{(j)}$ be the j th column of matrix A , then

$$A = [C^{(1)} \ \dots \ C^{(n)}]. \quad (1)$$

Now, consider the matrix multiplication in the following form

$$b = Ax = x_1C^{(1)} + \dots + x_nC^{(n)}, \quad (2)$$

which is a linear combination of A . □

1.2.4

Proof. Use contradiction, assume there is a nonzero y such that $Ay = 0$. Multiply A^{-1} on both sides, one has $y = A^{-1} \cdot 0 = 0$ which contradicts with $y \neq 0$. □

1.3.4

- Step 1

$$y_1 = 1.$$

- Step 2

$$y_2 = \frac{3 - (-1)y_1}{2} = 2.$$

- Step 3

$$y_3 = \frac{2 - 3y_1 - y_2}{-1} = 3.$$

- Step 4

$$y_4 = \frac{9 - 4y_1 - y_2 - (-3y_3)}{3} = 4.$$

Thus, $Y = [1 \ 2 \ 3 \ 4]^T$.

1.4.15

Proof. (a). For any $x \in \mathbb{R}^2$ and $x \neq 0$, consider $x^T A x = 4x_1^2 + 9x_2^2 > 0$, thus A is definite positive.

(b). Apply the method on the book, $R = \begin{bmatrix} 2 & 0 \\ 0 & 3 \end{bmatrix}$.

(c). For those diag elements of R , we do not force them to be positive square root, then the other three upper triangular matrices can be derived

$$R_2 = \begin{bmatrix} -2 & 0 \\ 0 & 3 \end{bmatrix}, \quad R_3 = \begin{bmatrix} -2 & 0 \\ 0 & -3 \end{bmatrix}, \quad R_4 = \begin{bmatrix} 2 & 0 \\ 0 & -3 \end{bmatrix}. \quad (3)$$

(d). Since for each r_{ii} , once it is determined, the i th row of R is uniquely determined. Thus, for an n by n matrix, there are two options for each row of matrix factor R , so 2^n possibilities totally.

□

1.5.9

Proof. Assume the width $s = k$ where $k \ll n$. Then, firstly consider the first k steps for solving from x_n to x_{n-k+1} , we have

$$x_i = (y_i - \sum_{j=i+1}^n r_{i,j} x_j) / r_{i,i}, \quad n - k + 1 \leq i \leq n. \quad (4)$$

Thus, the flops of first k steps is $f_1 = 1 + 3 + \dots + (2k - 1) = k^2$.

Next, consider the rest $n - k$ steps for solving from x_{n-k} to x_1 , since all $r_{i,j} = 0$ when $j - i > k$, we have

$$x_i = (y_i - \sum_{j=i+1}^{i+k} r_{i,j} x_j) / r_{i,i}, \quad 1 \leq i \leq n - k. \quad (5)$$

Thus, the flops of the rest steps is $f_2 = (n - k) \cdot (2k + 1)$.

Finally, the total flops is $f(n) = f_1 + f_2 = (n - k)(2k + 1) + k^2$. Consider $k \ll n$, $k^2 \approx 0$, and therefore $f(n) \approx (n - k)(2k + 1) \approx 2nk$.

□

1.7.10

(a). Use matlab command `det()` to find the determinants of four submatrice.

(b). Apply the Gaussian elimination for matrix A , the result upper triangular matrix

$$U = \begin{bmatrix} -6 & -1 & 2 & -3 \\ 0 & 2/3 & -1/3 & 2 \\ 0 & 0 & -1/2 & 3 \\ 0 & 0 & 0 & -3 \end{bmatrix}. \quad (6)$$

and

$$y = \begin{bmatrix} -14 \\ 8(1/3) \\ 10(1/2) \\ -12 \end{bmatrix}. \quad (7)$$

(c). Refer to 1.3.4, the solution is $x = [1 \ 2 \ 3 \ 4]^T$.